# HPC Programming

Trivial Parallelisation and OpenMP, Part I

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### Survey Outcome

- Programming languages: c, c++, python
- software used: Open source, but some are large (PandaROOT, BES3)
- Single-core (only one exception)
- Storage: GB to 10 TB (no problem, access pattern?)
- RAM: MB to 20 GB (1-3GB/core optimal)
- Runtime: seconds to hours (seconds to short: compare to set up costs of job)
- Number of program calls per analysis: 2-10k
- Boundaries: CPU, RAM (both great for HPC!), storage

# Trivial parallelisation

### **Trivial Parallelisation**

- Submit a single core job multiple times
- Quick and often only solution for large software blobs (large packages used in collaborations)
  - No principal difference compared to running on your desktop computer
- limits:
  - required RAM (3GB/core)
  - licensees (Mathematica, max 10 concurrent usages in university for such uses cases)
  - shared scratch in node (200GB)
  - parallel filesystem (loading at start, writing back results) max.  $\rightarrow$  10-100 jobs in parallel
- Hint: use job arrays
- Disadvantage: No control over speedup (this is gambling)

### Sample Submit Script

- 1. Define and reserve resources (nodes with RAM)
- 2. Once allocated, run the executables as defined or interactively

More examples <u>https://mogonwiki.zdv.uni-</u> <u>mainz.de/dokuwiki/slurm\_su</u> <u>bmit</u>

#!/bin/bash												
#SBATCH -o /home/pbotte/test/myjob.%j.%N.out												
#SBATCH -D /home/pbotte/t	test/											
#SBATCH -J MyJobName												
#SBATCH -A m2_him_exp	← account (NOT your account)											
#SBATCH -N 1	Request number of nodes											
<b>#SBATCH</b> partition=himst	ter2_exp 🗲 partition											
#SBATCHmem-per-cpu=1G												
#SBATCHmail-type=FAIL												
#SBATCHmail-user=pbott	te@uni-mainz.de											
<pre>#SBATCHtime=8:00:00</pre>	← wall time (>run time)											

```
module load gcc/6.3.0
echo TEST...
srun myExecutable
```

Submit with: sbatch submitScript.sh

### Batch System: SLURM

- Introduction and docu:
  - <u>https://mogonwiki.zdv.uni-mainz.de/dokuwiki/slurm\_submit</u>
  - <u>https://slurm.schedmd.com/tutorials.html</u>
- account to use: m2\_himkurs / m2\_him\_exp
- Reservation: himkurs / empty for himster2 usage
- Submit into partition: parallel / himster2\_exp
  - srun --pty -p parallel -A m2\_himkurs -N 1 --reservation kurstest bash -i
- Check what is running: squeue -u pbotte
  - 1184615\_79 parallel N203r001 pbotte R 1:00:40 52 z[0367-0386,0403-0413,0430-0450]
  - SSH login into your occupied nodes possible: eg ssh z0367

### Batch System: SLURM

- Submit script for later execution (batch mode)
  - sbatch --partition=himster2\_exp
- Create job allocation and start a shell to use it (interactive mode)
  - salloc -p parallel -N 1 --time=02:00:00 -A m2\_himkurs -- reservation=himkurs
- srun: Create a job allocation (if needed) and launch a job step (typically MPI job)
  - srun --pty -p parallel -N 1 --time=02:00:00 -A m2\_himkurs -reservation=himkurs bash -i
- sattach: Connect stdin/out/err for an existing job

## Support: Use of our HPC resources

- you plan to use HIMster2 or Mogon2 for your analysis?
- Feel free to ask for support
- We will discuss your problem and find some starting point for you.

# Introduction OpenMP



### Introduction OpenMP



- 1. Hardware Anatomy
- 2. Motivation
- 3. Programming and Execution Model
- 4. Work sharing directives
- 5. Data environment and combined constructs
- 6. Common pitfalls and good practice

### Anatomy of a ccNUMA

- Different parallel processing concepts: pipelining, vector computing, multicore, ...
- cache coherent Non-Uniform Memory Access (ccNUMA, AMD: 2003, industry wide: 2011)



• ccNUMA uses inter-processor communication between cache controllers

#### Machine (32GB)

Socket P#0 (16GB)								Socket P#1 (16GB)								
NUMANode P#0 (8192MB)								NUMANode P#2 (8192MB)								
L3 (8192KB)								L3 (8192KB	.3 (8192KB)							
L2 (2048KB)		L2 (2048KB)		L2 (2048KB)		L2 (2048KB)		L2 (2048KB	L2 (2048KB)		L2 (2048KB)		L2 (2048KB)		L2 (2048KB)	
L1i (64KB)		L1i (64KB)		L1i (64KB)		L1i (64KB)		L1i (64KB)		L1i (64KB)		L1i (64KB)		L1i (64KB)		
Lld (16KB)	L1d (16KB)	L1d (16KB)	L1d (16KB)	L1d (16KB)	Lld (16KB) Lld (16KB)		L1d (16KB)	L1d (16KB)	L1d (16KB)	L1d (16KB)	Lld (16KB)	L1d (16KB)	Lld (16KB)	L1d (16KB)	L1d (16KB)	
Core P#0 PU P#0	Core P#0         Core P#1         Core P#2         Core P#3         Core P#4         Core           PU P#0         PU P#1         PU P#2         PU P#3         PU P#4         PU		Core P#5 PU P#5	Core P#6 PU P#6	Core P#7 PU P#7	Core P#0 PU P#16	Core P#1 PU P#17	Core P#2 PU P#18	Core P#3 PU P#19	Core P#4 PU P#20	Core P#5 PU P#21	Core P#6 PU P#22	Core P#7 PU P#23			
NUMANode P#1 (8192MB)								NUMANode P#3 (8192MB)								
L3 (8192KB)								L3 (8192KB)								
L2 (2048KB)		L2 (2048KB)		L2 (2048KB)		L2 (2048KB)		L2 (2048KB)		L2 (2048KB)		L2 (2048KB)		L2 (2048KB)		
L1i (64KB)		L1i (64KB)		L1i (64KB)		L1i (64KB)		L1i (64KB)		L1i (64KB)		L1i (64KB)		L1i (64KB)		
L1d (16KB)	L1d (16KB)	L1d (16KB)	L1d (16KB)	L1d (16KB)	L1d (16KB)	L1d (16KB)	Lld (16KB)	L1d (16KB)	Lld (16KB)	L1d (16KB)	Lld (16KB)	L1d (16KB)	Lld (16KB)	L1d (16KB)	L1d (16KB)	
Core P#0 PU P#8	Core P#1 PU P#9	Core P#2 PU P#10	Core P#3 PU P#11	Core P#4 PU P#12	Core P#5 PU P#13	Core P#6 PU P#14	Core P#7 PU P#15	Core P#0 PU P#24	Core P#1 PU P#25	Core P#2 PU P#26	Core P#3 PU P#27	Core P#4 PU P#28	Core P#5 PU P#29	Core P#6 PU P#30	Core P#7 PU P#31	

• Output of hwloc tool: Topology of a ccNUMA Bulldozer server, 2 socket sytem

### Anatomy of a cluster computer

• N \* ccNUMA = cluster:



- Fast lossless interconnect: OmniPath between ccNUMA nodes
- Inside a node: NUMA, ccNUMA
- Multiple nodes: Distributed memory parallelisation (DMP)

### Anatomy of a cluster computer

• Latencies:



(all numbers are platform dependent)

### Introduction OpenMP



• Why OpenMP?

Relatively easy way: single-core program  $\rightarrow$  multicore shared-memory

- Released: 1997; widely and actively supported; currently version 4.5
- Only Fortran and C
- Overview:
  - OpenMP is a standard programming model for shared memory parallel programming
  - Set of compiler directives and a few library routines
    - efficient
    - → less problems during runtime (on ccNUMA nodes!) compared to library based shared RAM synchronisation (Pthreads)
  - **Portable**: Large set of compilers and hardware architectures
  - Slow start, direct results: step-by-step introduction of parallelisation
  - Shared memory: results in good speedup
- Prerequisite
  - A error free single-core program

### Implementation

Most modern compilers have support integrated, check their support

- Microsoft Visual C++ >2005,
- Intel Parallel Studio (OpenMP 3.1 since version 13),
- GCC ab Version 4.2 (OpenMP 4.0 since version 5.0),
- Clang/LLVM (OpenMP 3.1 since Version 3.6.1),
- Oracle Solaris Studio,
- Fortran, C und C++ Compiler der Portland Group (OpenMP 2.5),
- ...
- Homepage: openmp.org User group: compunity.org

## Comparison OpenMP / MPI

### OpenMP

- shared memory directives (compile time)
  - to define work decomposition
  - no data decomposition (data in shared memory)
- synchronisation is implicit

### Possible speedup:

- memory limited: Total bandwidth / single core bandwidth = 4 (hardware dependent)
- CPU limited: Number cores (+ possible cache effects)
- storage limited: do not use

## MPI (Message Passing Interface, later this course)

- software library (run time)
- user defines:
  - distribution of work & data
  - communication (when and how)

### Possible speedup:

- Per node limits: see OpenMP
- RAM/CPU limited: utilisation of N nodes
- Storage limited: ? (use node local scratch)

### Where to start?

Optimise your gain = speedup / work!

- 1. Try trivial parallelisation.
- 2. Parallelise your code with OpenMP, concentrate on time-consuming sections
- 3. Introduce MPI

large problems, work in team, check about available resources first (man power + hardware)

(different, if you join a group with existing MPI-code)

4. Hybrid programming: OpenMP + MPI to gain the last 10% speedup



### Glimpse: 1<sup>st</sup> OpenMP code

• OpenMP focusses on parallel loops with independent iterations

```
int main() {
    int in[100], out [100];
    for (int i=0; i<100; i++) {
        out[i] = MyLongFunc(in[i]);
        }
    }
}</pre>
int main() {
    int main() {
        int in[100], out [100];
        #pragma omp parallel for
        for (int i=0; i<100; i++) {
            out[i] = MyLongFunc(in[i]);
        }
}
```

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## **OpenMP: Programming Model**

- shared memory model
- Distribution of work between multiple threads ("workers")
  - Variables can be
    - Shared among all threads
    - Duplicated for each thread
  - Communication between threads through "barriers"
- Unintended data sharing  $\rightarrow$  race conditions (undefined behaviour) or dead lock
- To avoid race conditions: use synchronisation

read&write access to the same data by multiple threads

Due to different scheduling of threads between runs

### OpenMP: Execution Model



## **Execution Model Description**

- Begin execution as a single process (master thread)
- Fork-join of parallel execution
  - 1. Start of 1<sup>st</sup> parallel construct: Master thread creates N threads
  - 2. Completion of a parallel construct: threads synchronise (implicit barrier)
  - 3. Master thread continues execution
- At next parallel construct: work balancing with existing threads



### OpenMP Parallel Region Construct + Syntax

#pragma omp parallel [clause [, clause]]
 block

// emp end parallel

- *block* = to be executed by multiple threads in parallel. Each code executes the same code.
- Clause can be ("data scope"):
  - private (list) ← variables in list private to each thread & not initialised, standard for loop variables
  - shared (list) ← variables in list are shared among all thread, standard
  - firstprivate, lastprivate, threadprivate, copyin, reduction



### **Compiler Directives**

- #pragma directives #pragma omp directive\_name [clause [, clause]]
- Conditional compilation
   #ifdef \_OPENMP
   block, eg. printf("OpenMP sup.ted.");
   #endif
- Include file for library routines with compiler directives: #ifdef \_OPENMP #include <omp.h> #endif

#### • Why this (good practice)?

- Keep your code single-core and multi-core
- Do not copy your code (fork), always modify the main branch!

 $\rightarrow$  after years of development: main branch developed and your code is parallel but old!

### **Environment Variables**

- OMP\_NUM\_THREADS
  - Sets the number of threads
  - Set before execution, not during compilation
  - Bash: export OMP\_NUM\_THEADS=16 csh: setenv OMP\_NUM\_THEADS 16
- OMP\_SCHEDULE
  - Applies only to do/for and parallel do/for directives that have the schedule type RUNTIME
  - Sets schedule type and chunk size for all such loops
  - Bash: export OMP\_SCHEDULE="GUIDED,4" csh: setenv OMP\_SCHEDULE "GUIDED,4"

### **Runtime Library**

(Most of the OpenMP functionality arises from the compiler during compilation, but...)

- Query, runtime and lock functions comes from omp.h library #include <omp.h> (Implementation dependent)
- int omp\_get\_num\_threads()
  returns the current number of threads (N) executing the parallel region from which it is called
- int omp\_get\_thread\_num() return the thread number (0..N-1). Master thread is always 0
- wall clock timers: (similar to MPI\_WTIME in MPI) double omp\_get\_wtime(); provides elapsed time in a thread (needs not to globally consistent!)

```
# ifdef _OPENMP
double wt1,wt2;
wt1=omp_get_wtime();
# endif
//heavy computing
```

# ifdef \_OPENMP
wt2=omp\_get\_wtime();
printf( "wct %12.4g sec\n", wt2-wt1 );
# endif

### Exercise 1: Parallel region

Learning objectives:

- Runtime library calls
- Conditional compilation
- environment variables

#### Steps:

- 1. Copy the skeleton files from the course web page (see next slide)
- 2. Compile and run as serial program
- 3. Compile as openmp program (-fopenmp with cc) and run with different numbers of threads
- 4. Compare the run times between serial and openmp program

### Computational task

• Computational intensive function Karl Weierstraß; 1841:

$$\pi = \int_{-\infty}^\infty rac{\mathrm{d}x}{1+x^2} = 2 \cdot \int_{-1}^1 rac{\mathrm{d}x}{1+x^2}$$

• Run a analysis with several runs, do statistics

## Set up your workbench

- Connect 2 to Mogon2 / HIMster2 via SSH srun --pty -p parallel -N 1 --time=02:00:00 -A m2\_himkurs --reservation=himkurs -N 1 bash –I
  - 1. Use the first SSH connection for editing (gedit, vi, vim, nano, geany) and compiling: cc -fopenmp -o hello hello.c
  - 2. Use the second connection for the interactive execution on the nodes (no analysis on the head node!): OMP\_NUM\_THREADS=4 ./hello
- Download the files via: wget <u>https://www.hi-mainz.de/fileadmin/user\_upload/IT/lectures/WiSe2018/HPC/files/02.zip</u> && unzip 02.zip

#### Hints:

- Check compiler version: cc -V
- Run: OMP\_NUM\_THREADS=4 ./hello or export OMP\_NUM\_THREADS 4
- Possible to check reservation with: squeue -u USERNAME

### Exercise 2: Parallel region

Learning objectives:

• Parallel regions, private and shared clauses

Steps:

- 1. Use the code from exercise 1, and compile as openmp program (-fopenmp with cc) and run with number of threads=4
- Add a parallel region that prints the rank and the number of threads for each thread
   → expectation: undefined sequence of printf statements. No parallelisation of computation.
- 3. Try to create a race condition by:
  - First writing into registers: myrank = omp\_get\_thread\_num(); num\_threads = omp\_get\_num\_threads();
     And replace # pragma omp parallel private(myrank, num\_threads) →
     dependent result # pragma omp parallel
- 4. Add a #else directive that prints if the program was not compiled with OpenMP