HPC Programming

Message Passing Interface (MPI), Part IV

Peter-Bernd Otte, 11.12.2018

Introduction MPI

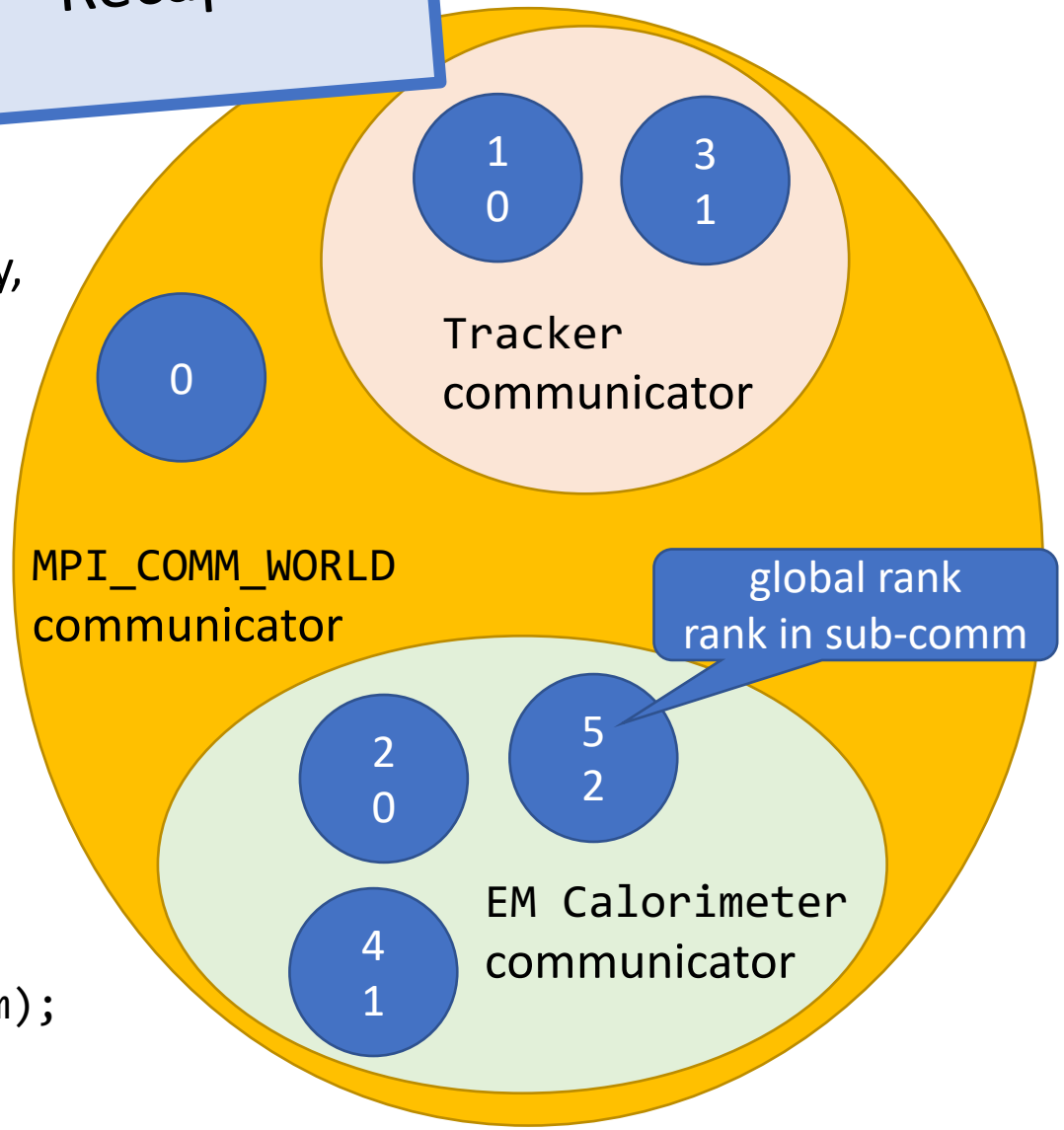
Recap

1. Overview / Getting Started
2. Messages & Point-to-point Communication
3. Nonblocking Communication
4. Error Handling
5. Groups & Communicators
6. Collective Communication
7. MPI I/O
8. MPI Derived Datatypes
9. Common pitfalls and good practice (“need for speed”)
10. Debugging and Profiling

MPI: MPI_Comm_split

Recap

- Creates new communicators based on colors
- int MPI_Comm_split(MPI_Comm comm, int color, int key, MPI_Comm *newcomm)
  - ordering in new group:
    - key == 0 → as sorted in old
    - key != 0 → according to key values
  - one member group: color = MPI_UNDEFINED



- Example:

```
MPI_Comm newcomm;
MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);
mycolor = my_rank/3;
MPI_Comm_split(MPI_COMM_World, mycolor, 0, &newcomm);
MPI_Comm_rank(newcomm, &my_new_rank);
```

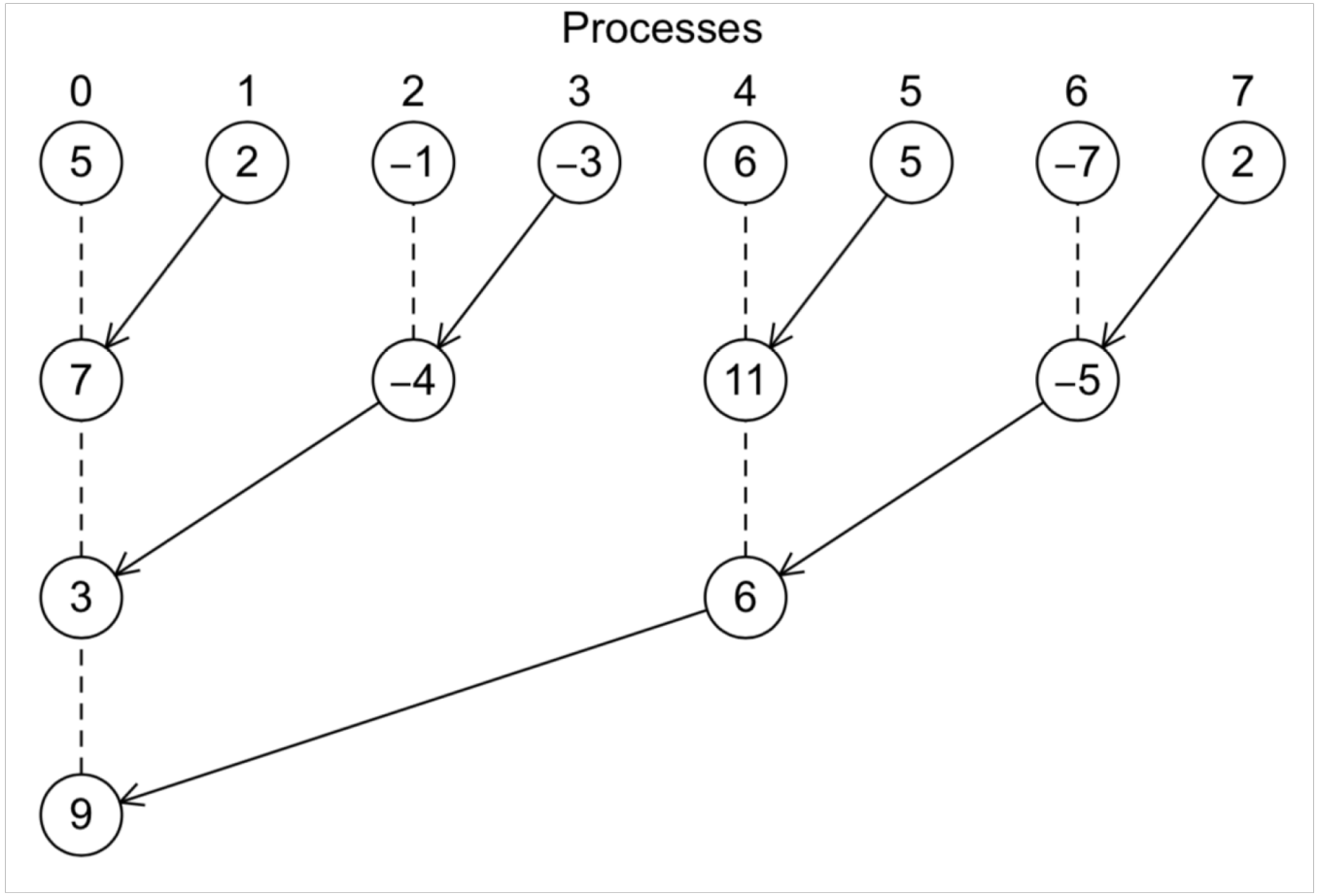
MPI: MPI_Reduce

Recap

- Reduces values on all processes to a single value
  (eg global sum)

```
int MPI_Reduce(
void *sendbuf /*in*/,
void *recvbuf /*out*/,
int count /*in*/,
MPI_Datatype datatype /*in*/,
MPI_Op operator /*in*/,
int dest_process /*in*/,
MPI_Comm comm /*in*/)
```

- hints:
  - with count>1, MPI can operate on arrays
  - sendbuf and recvbuf need to different (no aliasing!)



MPI: P2P ⇔ Collective Communicatio

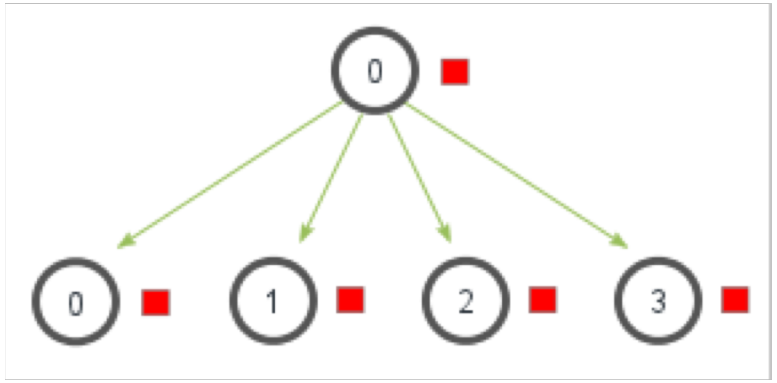
Recap

- ALL processes in communicator must call SAME collective function at the same time.
- Arguments in all ranks must fit:
  - eg. same dest_process, datatype, operator, comm
  - depending on function
- Only rank dest_process may use recvbuf (but all ranks have to provide such argument)
- MPI_Reduce calls matched solely on:
  - the communicator and
  - the order on which they are called.
  - No helping tags or sender id available.

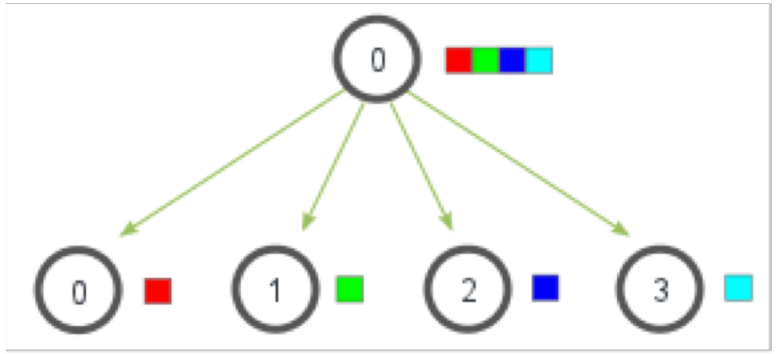
MPI: Broadcast and Scatter

Recap

Broadcasts the same message from the process ”sending_rank” to all other processes of the communicator



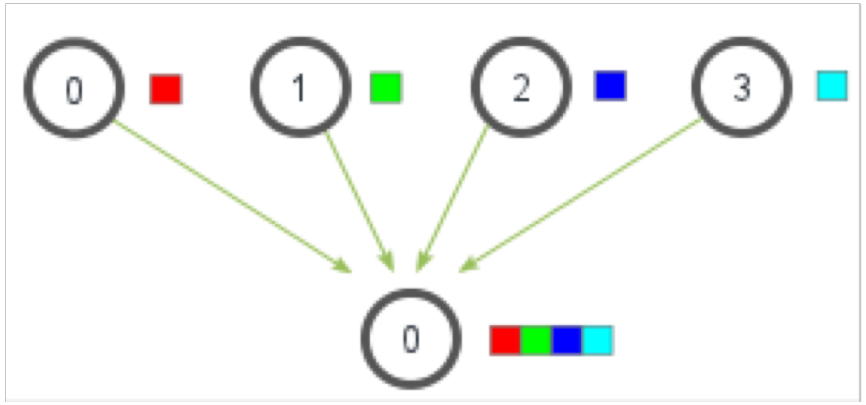
Scatter: Sends data from one process to all other processes in a communicator



MPI: MPI_Gather

Recap

Gathers together values from a group of processes



- MPI_Gather(
  void* send_data /*in*/,
  int send_count /*in*/,
  MPI_Datatype send_datatype /*in*/,
  void* recv_data /*out*/,
  int recv_count /*in*/,
  MPI_Datatype recv_datatype /*in*/,
  int dest_proc /*in*/,
  MPI_Comm comm /*in*/)

- Special cases: MPI_Gatherv

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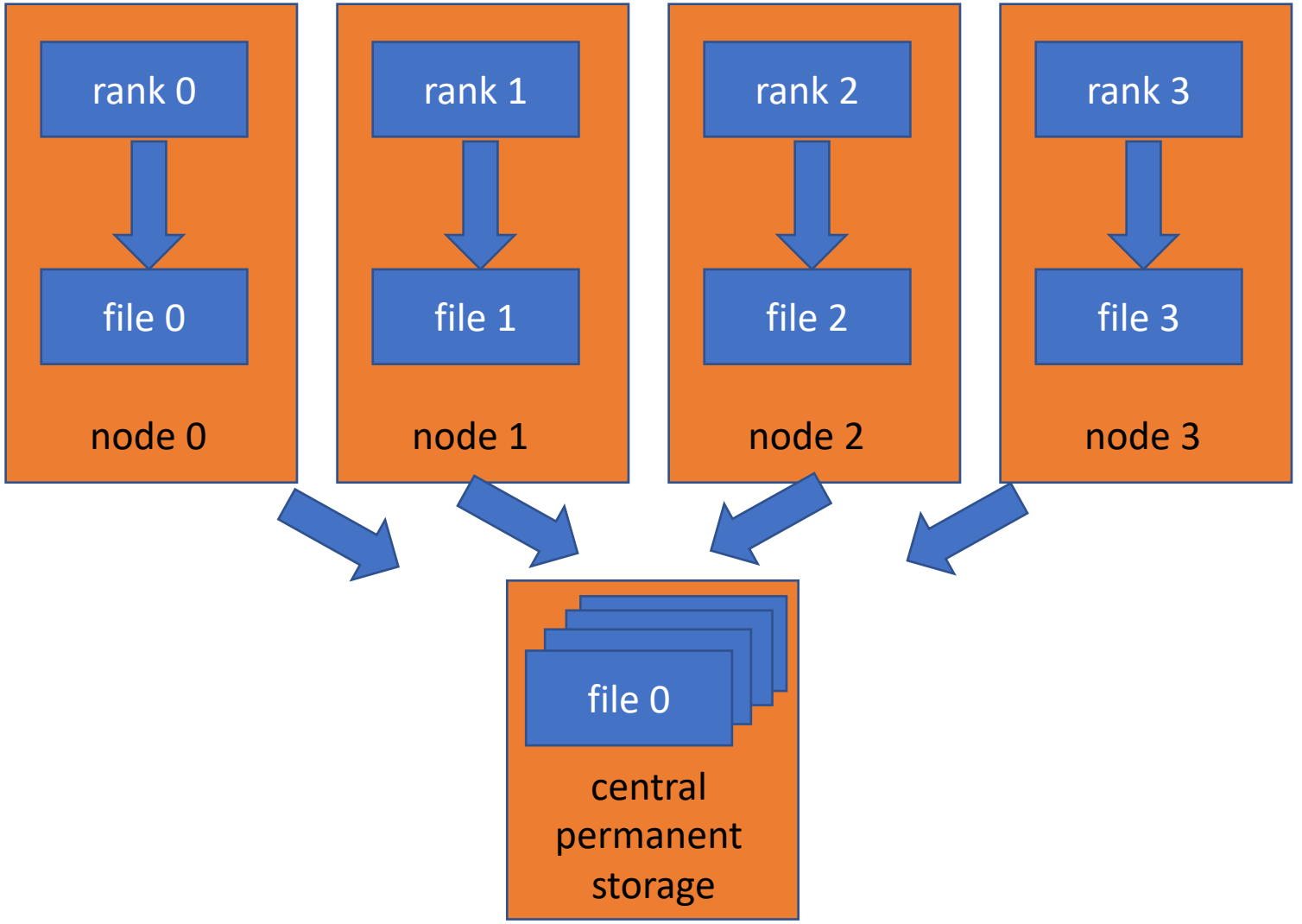
Motivation: MPI I/O 1

- Standard (POSIX): each process writes to a single separate file on scratch(!) device
- Typical situation: analysis framework
- parallel → scales!

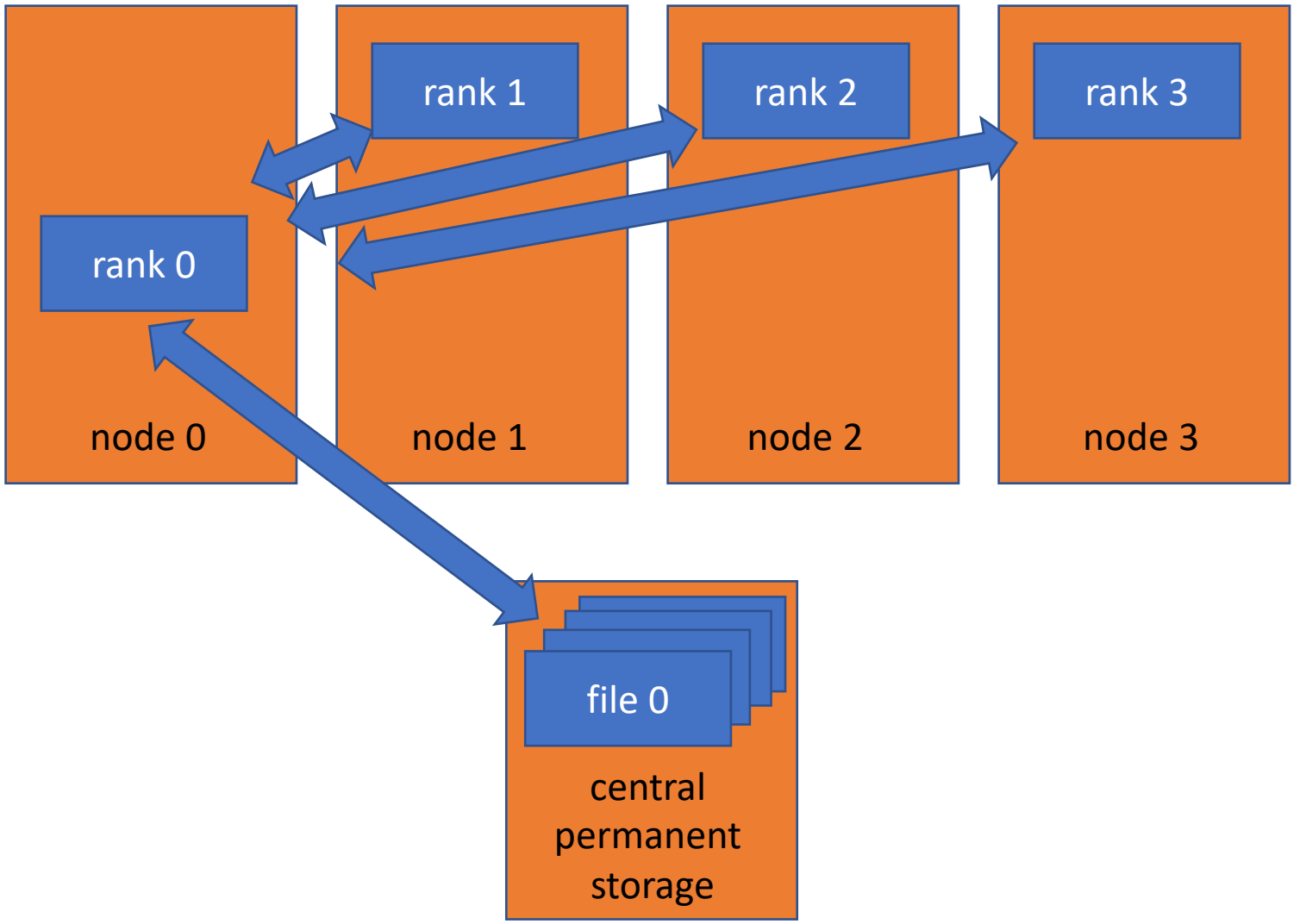
- collection of all these single files → serialisation or worse
- many files → bad for meta data server



Motivation: MPI I/O 2

- Legacy: only single rank reads/writes
- Typical situation: apps recently parallelised, OpenQCD

- serial access and broadcast (→ worse than
- reads only a fraction of a file → bad for meta data server

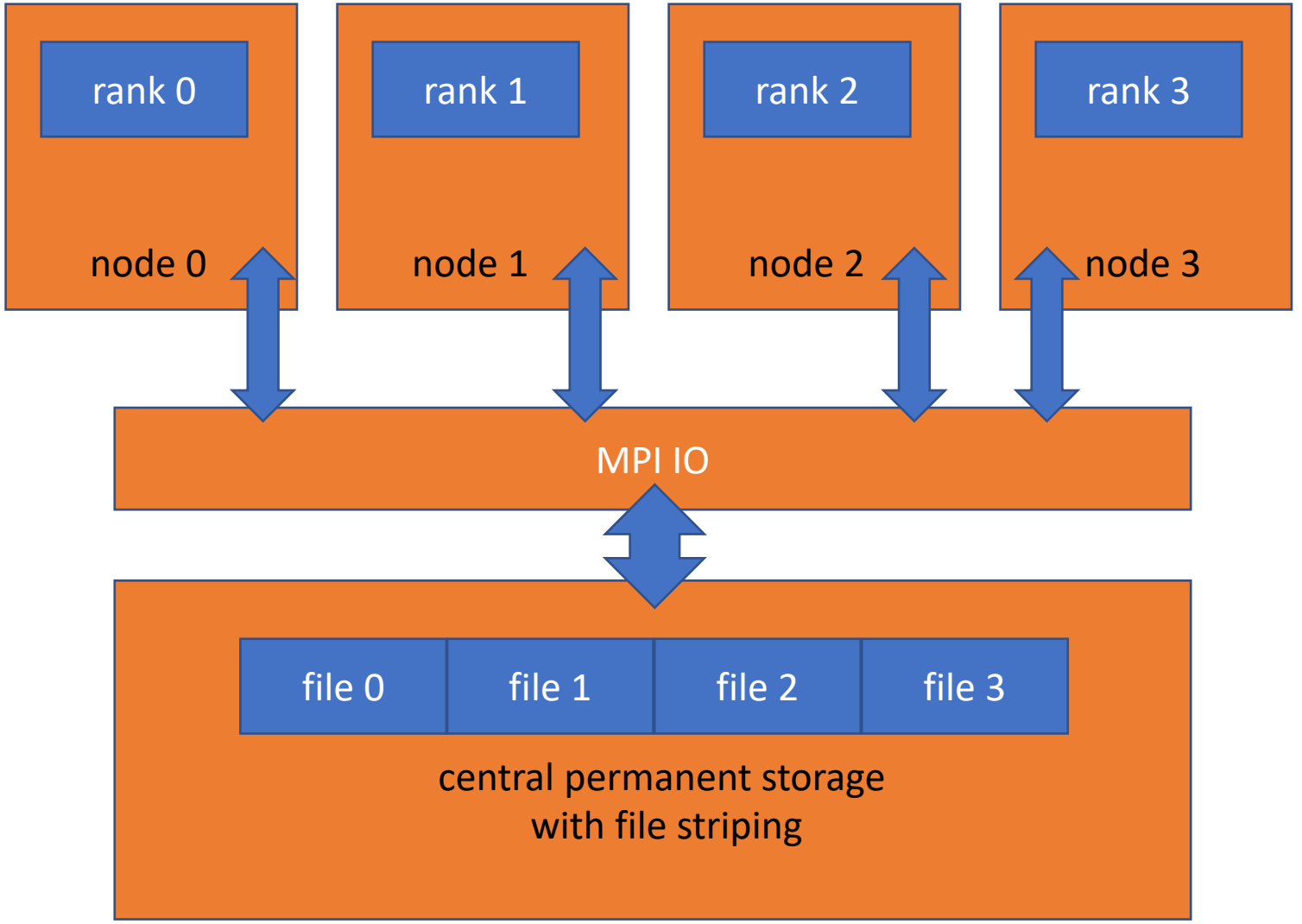


Motivation: MPI I/O 3

- Speed up with cooperation and parallelism

MPI IO:

- simultaneous access cooperation
- single file
- provides replacement function for POSIX



Motivation: MPI & MPI IO 4

MPI I/O is based on:

- MPI & parallel FS (→ fast)
- handle read/write accesses like sending/receiving of messages

parallel I/O requirements	analogy on MPI
collective file operations	MPI communicators
non-contiguous access	MPI derived datatypes
nonblocking operations	MPI functions with immediate return in combination with Wait.

not yet discussed in this lecture

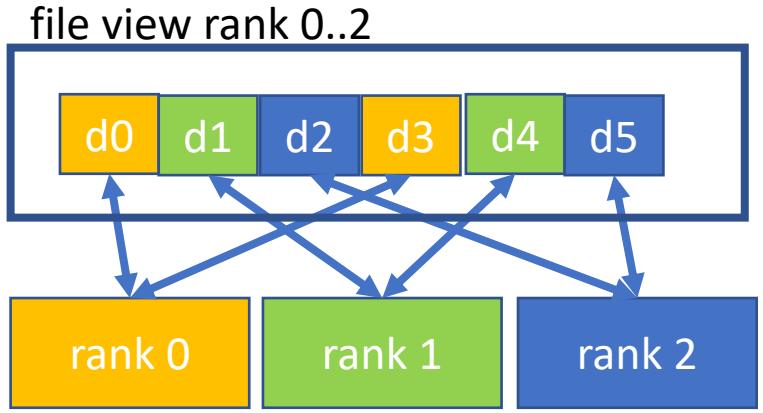
MPI IO principles

- MPI file contains elements of a single MPI datatype (“etype”)
- rank file access provided by access templates
- read/write routines in MPI IO: nonblocking / blocking and collective / individual reads
- file pointers: individual and shared
- automatic data conversion in heterogenous systems

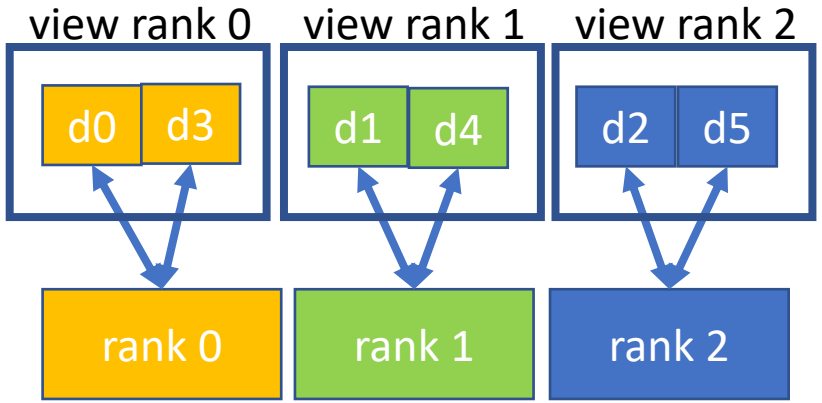
MPI: Access possibilities

- Array of data in file

- 3 ranks processing this file
  1. full view on file for every rank (like standard POSIX) with MPI_File_write_at()



  2. reduced view on file for every rank with MPI_File_set_view() and MPI_File_write()



MPI IO: Opening a file

- int MPI_File_open (
  MPI_Comm comm,
  ROMIO_CONST char *filename,
  int amode,
  MPI_Info info,
  MPI_File *fh /*out*/)
- collective within communicator.
  - all processes in comm. call function with same arguments (filename, amode)
  - process-local files with MPI_COMM_SELF as communicator
- returns a file handle
  - representing the file, communicator and the current view (see next slides)
- default:
  - displacement = 0, etype=MPI_BYTE → each process has access to whole file (”slide before: full view”)
- No info = MPI_INFO_NULL, otherwise provide timeouts, buffer sizes or stripe factors here.

MPI IO: Access Mode

- remember: same amode argument on all processes (collective!)
- combe these arguments bit wise → Operator | (better not +)
- Be as restrictive as possible to allow for storage optimisation

Constants	
MPI_MODE_APPEND	all file pointers set to end of file
MPI_MODE_CREATE	Create the file if it does not exist.
MPI_MODE_DELETE_ON_CLOSE	
MPI_MODE_EXCL	Error creating a file that already exists.
MPI_MODE_RDONLY	Read only.
MPI_MODE_RDWR	Reading and writing.
MPI_MODE_SEQUENTIAL	only sequential access, eg: tapes
MPI_MODE_WRONLY	Write only.
MPI_MODE_UNIQUE_OPEN	file not opened concurrently

caution: any following call of MPI_FILE_SET_VIEW will reset this to 0

MPI IO: Closing a file

- collective function
- int MPI_File_close(MPI_File *fh)

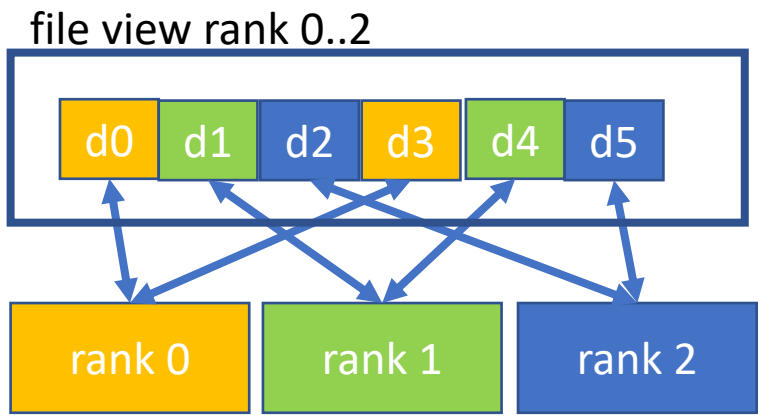
MPI IO: File Deletion

1. int MPI_File_delete(
   ROMIO_CONST char *filename,
   MPI_Info info)
   - file need not be currently opened

2. Provide argument „amode = MPI_MODE_DELETE_ON_CLOSE“ in MPI_File_Open

MPI IO: Writing to file with explicit offset

- (needed for exercise 7)
- int MPI_File_write_at(
  MPI_File fh,
  MPI_Offset offset,
  ROMIO_CONST void *buf,
  int count,
  MPI_Datatype datatype,
  MPI_Status *status)
  - buffer includes min count elements of type datatype



- writes count times elements from buffer to to the file
- starting at offset * sizeof(datatype) from begin of view

MPI IO: Reading from a file with explicit offsets

- int MPI_File_read_at(
  MPI_File fh,
  MPI_Offset offset,
  void *buf,
  int count,
  MPI_Datatype datatype,
  MPI_Status *status)
- read count elements of datatype
- starting at offset * sizeof(datatype) from begin of view
- EOF is reached, once amount of data read < count
  - use MPI_Get_Count(status, datatype, received_count)
  - note: EOF is no error

MPI IO: Individual file pointers 1/2

- int MPI_File_read(MPI_File fh, void *buf, int count, MPI_Datatype datatype, MPI_Status *status)
- int MPI_File_write(MPI_File fh, ROMIO_CONST void *buf, int count, MPI_Datatype datatype, MPI_Status *status)
- same functions as those functions with “_at”, except:
  - each process has it’s private current value of file offset (“file pointer”)
  - after access, private offset updates:
    - private offset points to the next datatype of the last accessed.

MPI IO: Individual file pointers 2/2

- int MPI_File_seek(MPI_File fh, MPI_Offset offset_new, int whence /*Update mode*/)
  - Update mode = MPI_SEEK_SET → set private file offset to offset_new
  - MPI_SEEK_CUR → advance private file offset by offset_new
  - MPI_SEEK_EOF → set private file offset to EOF + offset_new

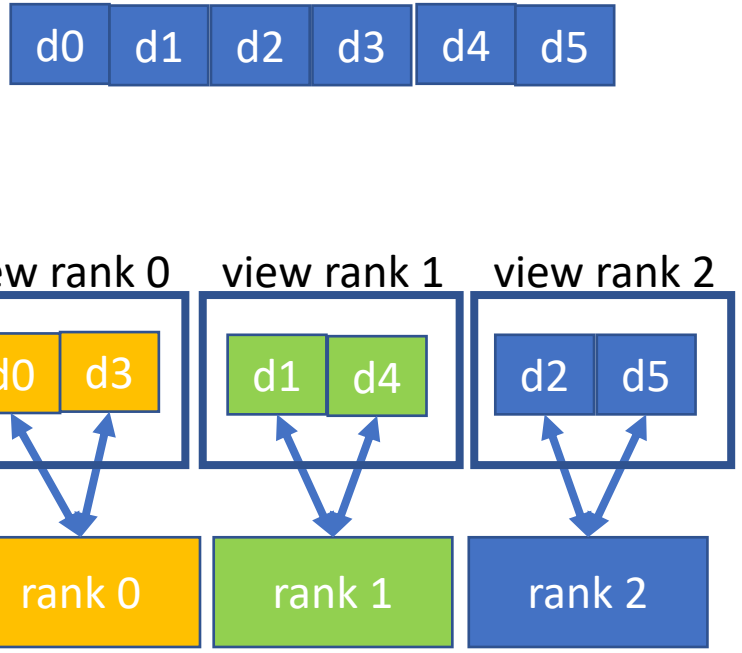
inquire offset:

- int MPI_File_get_position(MPI_File fh, MPI_Offset *offset)
- int MPI_File_get_byte_offset(MPI_File fh, MPI_Offset offset, MPI_Offset *disp)
  - disp = absolute byte position of offset (nonnegative integer)

- To convert an offset into byte displacement (needed eg for a new view)

MPI IO: File views 1/2

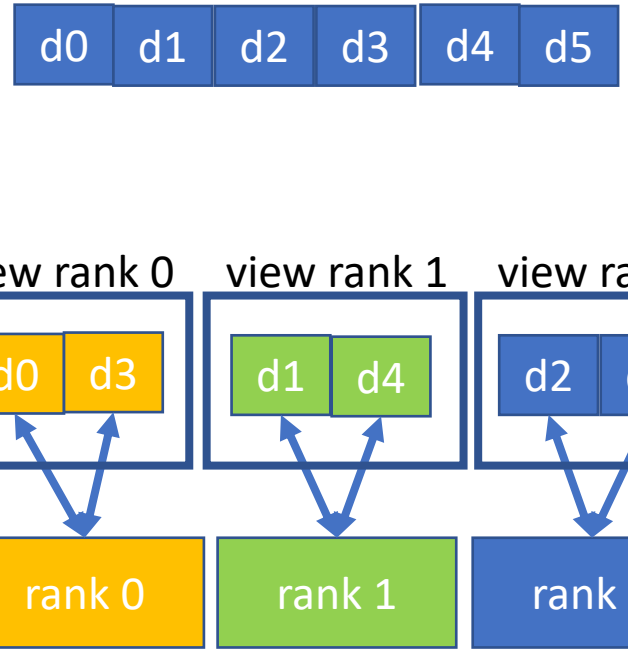
- Each process gets a separate view of the file, collective operation (necessary for exercise 8)
- Defined by (Displacement, datatype, filetype)
  - Standard = (0, MPI_BYTE, MPI_BYTE) = linear byte stream
- can be changed during runtime
- int MPI_File_set_view(
  MPI_File fh,
  MPI_Offset disp,
  MPI_Datatype etype,
  MPI_Datatype filetype,
  ROMIO_CONST char *datarep /*see next slide*/,
  MPI_Info info)
- Get view via MPI_File_get_view()



MPI IO: File views 2/2

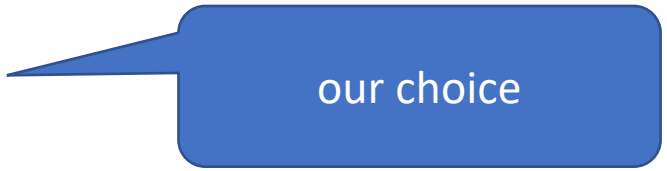
- Worked out example, create MPI_Type filetype first:

```
etype = MPI_CHAR;
ndims = 1;  /*dimensions of following arrays*/
array_of_sizes[0]    = 3;
array_of_subsizes[0] = 1;
array_of_starts[0]   = my_rank;
MPI_Type_create_subarray(ndims,
     array_of_sizes, array_of_subsizes,
     array_of_starts, MPI_ORDER_C, etype,
     &filetype);
MPI_Type_commit(&filetype);
MPI_File_set_view(fh, 0, etype, filetype,…);
```



MPI IO: Data representation

- native:
  - data in file = data in memory
  - no type conversions (no loss of precision and I/O performance) on homogenous systems
  - not possible on heterogenous systems
  - no guarantee by MPI to mix C and Fortran
- internal:
  - implementation dependent, for heterogenous systems
- external32
  - follows standardized representation (IEEE)
  - all input/output according to “external32” representation → interoperable between different MPI impl.
  - due to type conversions from/to native: data precision and I/O performance is reduced
  - can be read/written also by non-MPI programs



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Most common MPI pitfalls:

- FIRST: optimise single core performance
- efficiency of MPI application-programming is not portable
  → optimize for every system needed (when aiming for highest speeds)

- Most Common pitfalls:
  - Deadlocks and serialization
  - Late sender
  - Late sender

- further hints:
  - Overlap communication and computation
  - Global communication involving many or all MPI processes include costly synchronizations.
    - combine such reductions to overhead
  - try to share huge buffers instead of copying
  - Check resources, try to avoid local swap → use more machines, less ranks / threads per node

MPI optimisation

- Advanced:
  - Contention:
    - Miss ratio senders / receiver,
    - low bisectional bandwidth between nodes,
    - non ideal network routing
  - Non optimal domain decomposition (slicing your detector, try slices with smaller surfaces):
    - Try different “data decomposition (divide the problem differently)”
    - too much communication overhead,
    - as many ranks on a single node → avoid network
  - On multi socket systems: sending rank should be on core in hardware, which is closest to network link
  - Check for load imbalances, use tuning tools

MPI optimisation: Binding

- Binding processes and their threads prevents the OS scheduler from moving them across the available CPU sockets or cores.
- **Memory-bound MPI application** with one MPI process per socket
  - $MPIEXEC $FLAGS_MPI_BATCH --map-by ppr:1:socket --bind-to core a.out
- **Compute-bound MPI application** with as many processes per node as there are cores
  - $MPIEXEC $FLAGS_MPI_BATCH --bind-to core --map-by core a.out
- **MPI application with n processes per socket** (n < #cores)
  - For certain MPI applications that are neither completely compute- nor completely memory-bound it might be beneficial to run them with less processes per socket than cores are there.
  - $MPIEXEC $FLAGS_MPI_BATCH --map-by ppr:2:socket --bind-to core a.out
  - ## number of processes per socket ---^
- **Examining the Binding**, OpenMP: --report-bindings
- With OpenMP: $MPIEXEC $FLAGS_MPI_BATCH -x OMP_NUM_THREADS -x OMP_PLACES -x OMP_PROC_BIND -x KMP_AFFINITY \

  --map-by ppr:1:socket --bind-to socket a.out
- **Depends on MPI Implementation** (Intel: pnning) and if OpenMP is used.

MPI optimisation: Binding

```
mpirun -n 4 --report-bindings -bind-to core ring-sub
[login22:17413] MCW rank 0 bound to socket 0[core 0[hwt 0-1]]: [BB/../../../../../../../../..][../../../../../../../../../..]
[login22:17413] MCW rank 1 bound to socket 1[core 10[hwt 0-1]]: [../../../../../../../../../..][BB/../../../../../../../../..]
[login22:17413] MCW rank 2 bound to socket 0[core 1[hwt 0-1]]: [../BB/../../../../../../../..][../../../../../../../../../..]
[login22:17413] MCW rank 3 bound to socket 1[core 11[hwt 0-1]]: [../../../../../../../../../..][../BB/../../../../../../../..]
```

MPI: Possible sources of errors

1. Starting multi-core program: do not copy / fork your code, improve existing.
2. Error free single core program.
3. Hardware (CPU, RAM, network, storage) free of errors.

- program hangs
  send / receive do not match (sender it, communicator, tag, etc.) → verify parameters
- MPI_Send crashes:
  Buffer address correct? Still correct? eg OpenMP task gets executed with delay (use “omp taskwait”)
- MPI_Recv crashes: MPI library tells, msg is larger than recv buffer
  message from correct sender received? Did tags match? wrong message order? → use unique tag
- received message data is wrong
  Send buffer has been modified (buffered send) before sent / Received buffer has been accessed before arrival of data
- Using OpenMP and MPI in parallel:
  → Tell mpirun about it, use correct MPI multi-thread level (eg MPI_THREAD_SERIALIZED or MPI_THREAD_MULTIPLE)

MPI I/O

- Best practices of using MPI I/O:
  - make as few file I/O calls in general
  - in order to create big data requests and
  - have as few meta-data accesses (seeks, query or changing of file-size).

- Change MPI_Info key-values, according to your needs, eg:
  - ```
    MPI_Info info;
    MPI_Info_create(&info);
    /* Enable ROMIO's collective buffering */
    MPI_Info_set(info, "romio_cb_read", "enable");
    MPI_Info_set(info, "romio_cb_write", "enable");
    MPI_File_open (MPI_COMM_WORLD, fn, MPI_MODE_CREATE | MPI_MODE_WRONLY, info, &fh);
    ```
    ```

General File Access Hints

- Bad I/O performance due to:
  - Accessing that same portion of the file → locks
  - Other i/o in parallel
  - random accesses
  - datasize(i/o requests) << filesystem block size
  - files too small / too many files / too many open&closes → metadata servers overloaded
- Avoid data access:
  - Recalculate when it's faster
  - group small operations to larger chunks
  - Reduce data accuracy, possible? → less data!
- Helpful:
  - Use parallel I/O libraries: MPI I/O, HDF5, etc. and use their non-blocking MPI I/O routines
  - large and contiguous requests
  - Use derived datatypes to support MPI I/O in its work
  - Open files in the correct mode (eg only readonly) to allow for optimisations
  - Not too many open files at the same time
  - flushes only when absolutely necessary.
  - Create files independent of the number or processes (easier post processing and restarts with different rank size)

Optimisation

- Good read for further studies:
  - Hager, Wellein: “Introduction to High Performance Computing for Scientists and Engineers”, CRC Press

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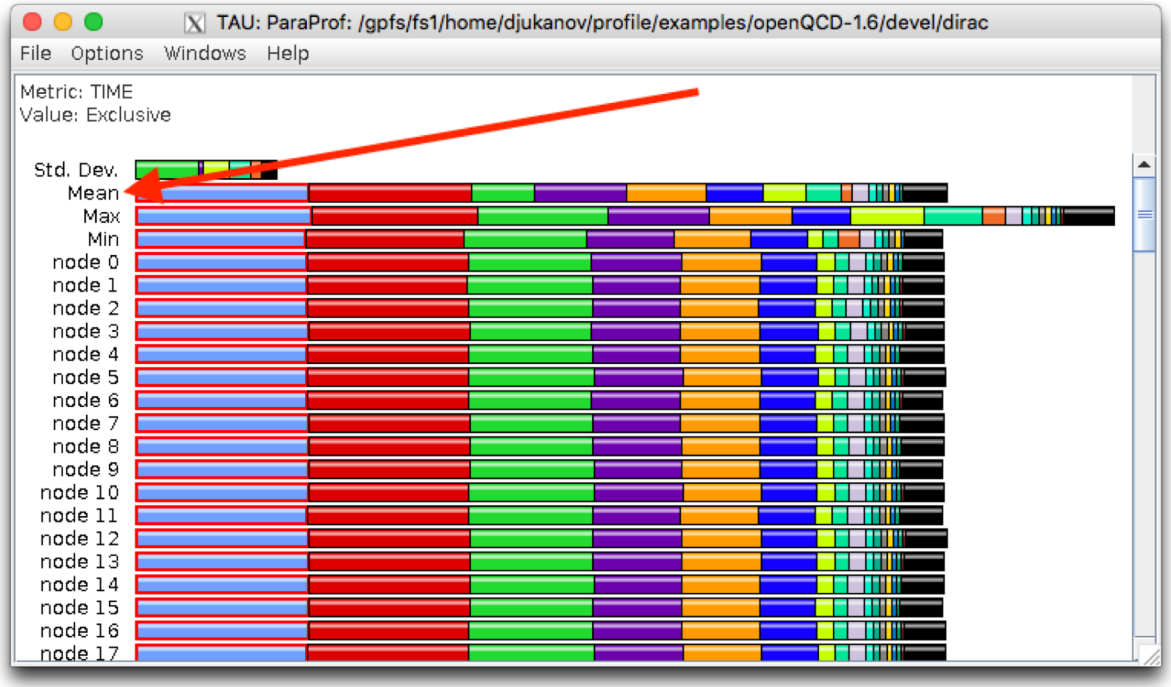
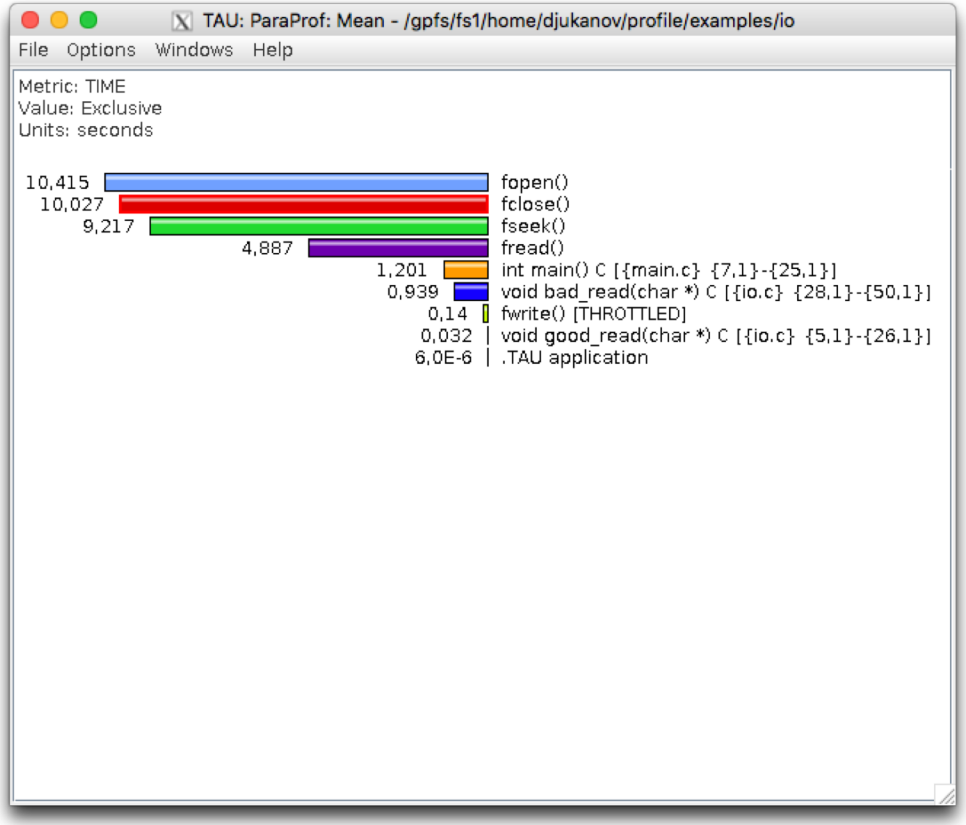
topics of future lectures

MPI: Profiling Glimpse

- See where time is spent
- Identify idle periods





Set up your workbench

- Connect 2 times via SSH to Mogon2 / HIMster2
  1. Use the first SSH connection for editing (gedit, vi, vim, nano, geany) and
     module load mpi/OpenMPI/3.1.1-GCC-7.3.0
     compiling: mpicc -o ExecutableName SourceFileName.c
  2. Use the second connection for the interactive execution on the nodes (no execution on the head node!):
     salloc -p parallel -N 1 --time=01:30:00 -A m2_himkurs --reservation=himkurs -C skylake
     module load mpi/OpenMPI/3.1.1-GCC-7.3.0
     mpirun -n 2 ./ExecutableName
- Download the files via: wget https://www.hi-mainz.de/fileadmin/user_upload/IT/lectures/WiSe2018/HPC/files/MPI-03.zip && unzip MPI-03.zip

Hints:

- If the reservation with salloc –p parallel fails, try:
  - salloc -p devel -n 4 -A m2_him_exp
- The reserved resources with salloc can’t be overwritten with mpirun
  - Resources(salloc) => Resources(mpirun)
- Possible to check reservation with: squeue -u USERNAME

Exercise 7:

Learning objectives:

- first usage of MPI IO and MPI_File_write_at()

Steps:

1. Download the skeleton from lecture webpage:
   - wget https://www.hi-mainz.de/fileadmin/user_upload/IT/lectures/WiSe2018/HPC/files/MPI-07.zip && unzip MPI-07.zip
2. Each rank writes 5 times its rank number into a common file (do not use more than 9 ranks). The output should look like (with 4 ranks):
   01230123012301230123

Hints:

- offset = my_rank + Comm_Size * i, i=0..4
- Each process uses the default view
- To write numbers as ASCII characters use
  buf = '0' + (char)my_rank;
- You can use “cat FILENAME” to check your written output.
- Real world hint: Your home directory is not a parallel FS. For full speed use /lustre/…

Exercise 8:

Learning objectives:

- Write to a file with MPI_File_set_view

Steps:

1. Download the skeleton from lecture webpage:
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2. Achieve the same result as in exercise 7 but make use of MPI_Type_create_subarray, MPI_File_set_view and MPI_File_write