

HPC Programming

Message Passing Interface (MPI), Part III

Peter-Bernd Otte, 26.11.2019

Introduction MPI

1. Overview / Getting Started
2. Messages & Point-to-point Communication
3. Nonblocking Communication
4. Error Handling
5. Groups & Communicators
6. Collective Communication
7. Dealing with I/O
8. MPI Derived Datatypes
9. Common pitfalls and good practice (“need for speed”)



MPI: different communications modes

Recap

	Blocking	Non-Blocking	note
standard send	MPI_Send	MPI_!Send	synchronous or asynchronous send (depending on message size and implementation) uses internal buffer.
synchronous send	MPI_SSend	MPI_!SSend	Only completes when the receive has started
asynchronous (buffered) send	MPI_BSend	MPI_!BSend	Completes after buffer copy (always).
ready send	MPI_RSend	MPI_!RSend	problematic: mandatory to have matching receive already listening. Not discussed in this lecture. Might be fastest solution.

„i” stands for immediate return

	Blocking	Non-Blocking	note
standard receive	MPI_Recv	MPI_IRecv	works for all sending routines.

MPI: P2P communications, Pros and Cons

- **synchronous send**
 - **risk of serialisation, waiting and/or deadlock**
 - high latency but best bandwidth
- **asynchronous send**
 - **no risks** (except: take care of your buffers)
 - low latency but bad bandwidth
- **standard send**
 - **risk of implementation and message dependence behaviour**
 - **plus risks of synchronous send**

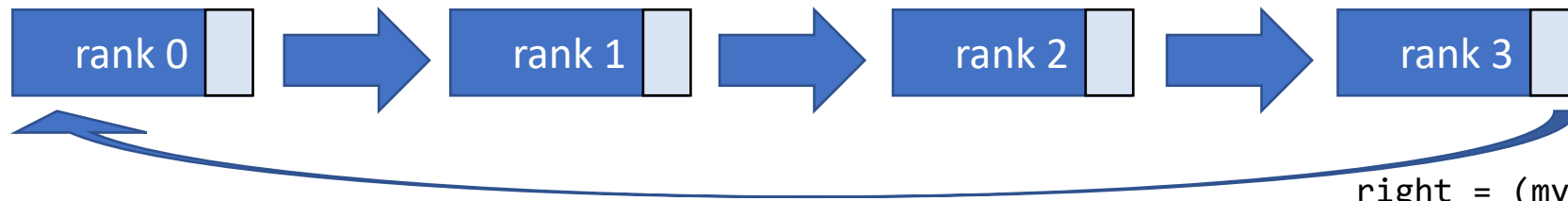


Recap

MPI: Non-Blocking Send & Receive

Recap

- to a 1D ring with 1 piece of data passing in one direction



```
right = (my_rank+1) % size;  
left = (my_rank-1+size) % size;
```

- **cyclic**: MPI_Send(...to right...)
MPI_Recv(...from left...)

deadlock!

All are waiting
for a receiver

- **non-cyclic**: for rank < size-2: MPI_Send(...to right...)
for rank > 0: MPI_Recv(...from left...)

serialisation!

highest rank starts,
rank 0 last

(hint: all this only true if MPI calls are synchronous sends)

MPI: Non-Blocking communication

Recap

This can be accomplished by:

- non-blocking send
 1. `MPI_Isend();`
 2. `Different_Work();`
 3. `MPI_Wait();` //Waits until `MPI_Isend` completed / send buffer is read out
- non-blocking receive
 1. `MPI_Irecv();`
 2. `Different_Work();`
 3. `MPI_Wait();` //Waits until `MPI_Irecv` completed / receive buffer is filled

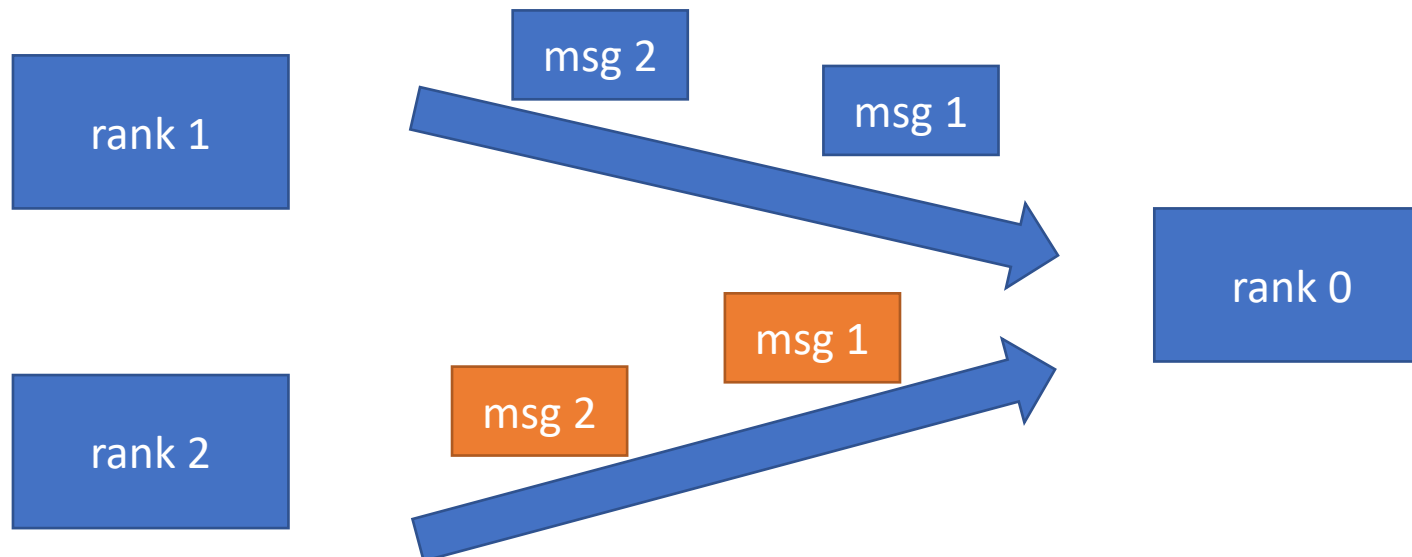
Golden MPI rule:
always ≤ 3 lines of
MPI_* calls per task

otherwise:
check MPI reference or
wrong coding

MPI: Message Order Preservation

Recap

- **Messages do not overtake**, if same:
 - communicator (eg MPI_COMM_WORLD),
 - source rank and
 - destination rank
- true for: synchronous and asynchronous communications
- messages from different senders can overtake





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MPI: Error handling

- in short, standard behaviour:
 - MPI: abort on error
 - MPI-IO: continue and just report
 - only if error is detected by MPI, otherwise unpredictable behaviour
- in detail:
 - most important foundation: hardware error free
 - CPU, RAM & network
 - have different techniques to detect hardware errors (eg ECC-RAM, checksums in network packages)
 - you (or your system admin) are informed if hardware problem occurs
 - Change standard behaviour:
`int MPI_Comm_set_errhandler(MPI_Comm comm, MPI_Errhandler errhandler)`
`int MPI_File_set_errhandler(MPI_File file, MPI_Errhandler errhandler)`

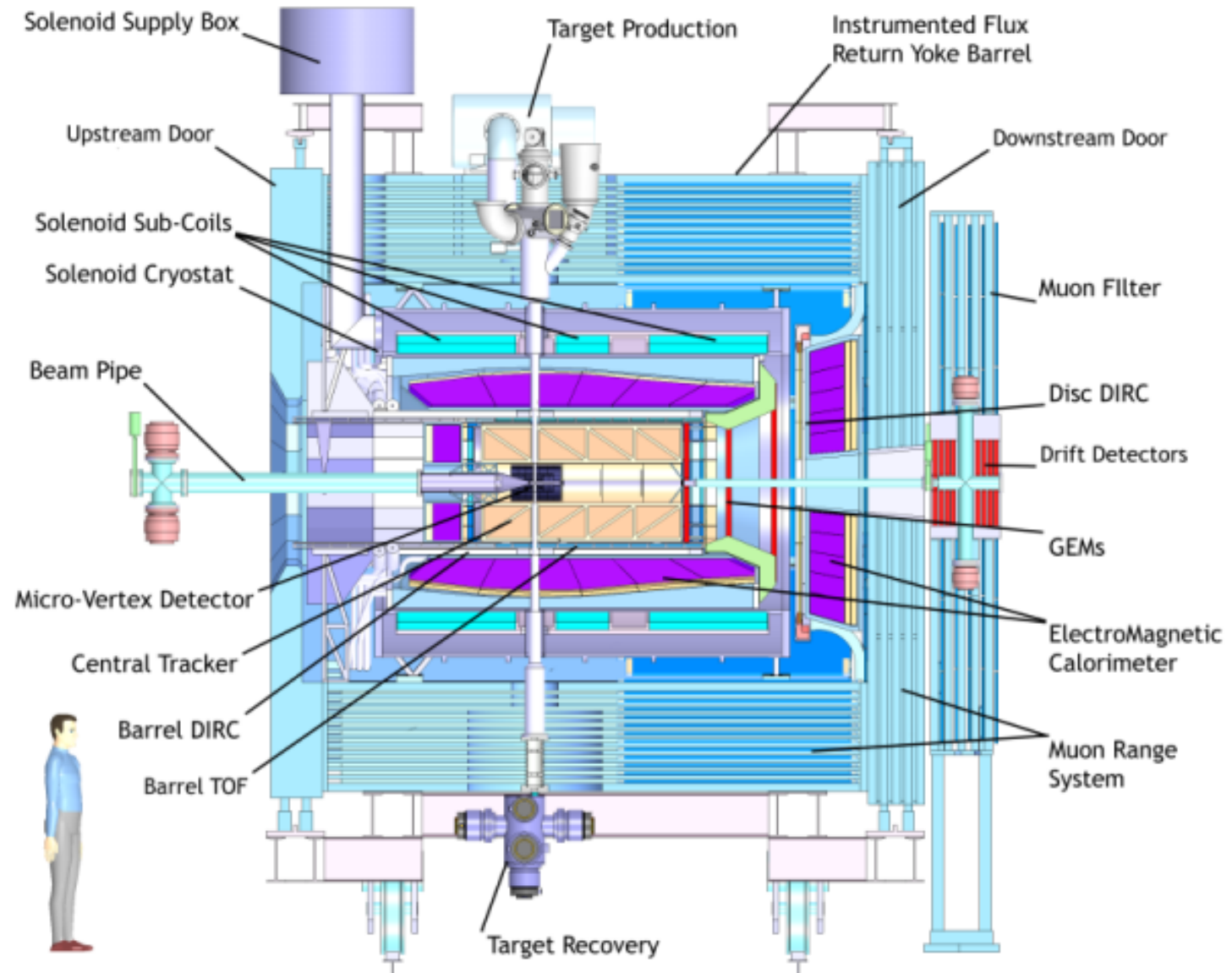


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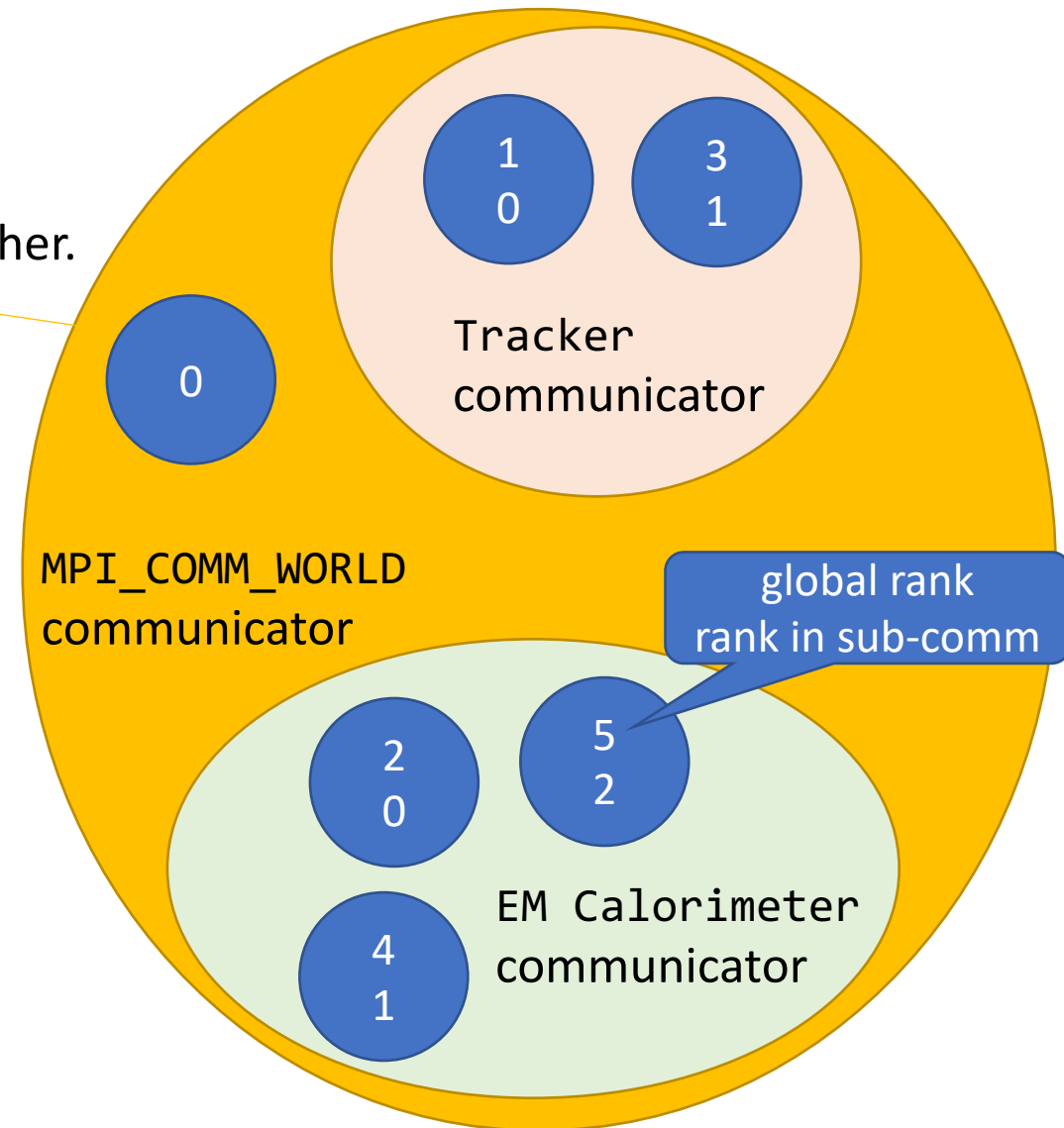
Motivation: Sub-Communicators

- Particle reconstruction
- Multiple layers in detector
- Multiple ranks working in several groups
 - code readability
 - collective communication within group
- OR: a library should NEVER use `MPI_COMM_WORLD` to not mix up with the main program.
- See Exercise 5



MPI: Sub-Communicators

- MPI Communicator = group of processes that can send messages to each other.
- All processes are in MPI_COMM_WORLD communicator
- Defining sub groups (eg readability, library):
 1. MPI_Comm_split
 2. MPI_Comm_group + MPI_Comm_create
- Number of members and size in communicator: MPI_Comm_size, MPI_Comm_rank



MPI: MPI_Comm_split

- Creates new communicators based on colors
- `int MPI_Comm_split(MPI_Comm comm, int color, int key, MPI_Comm *newcomm)`
 - ordering in new group:
 - `key == 0` → as sorted in old
 - `key != 0` → according to key values
 - one member group: `color = MPI_UNDEFINED`

- Example:

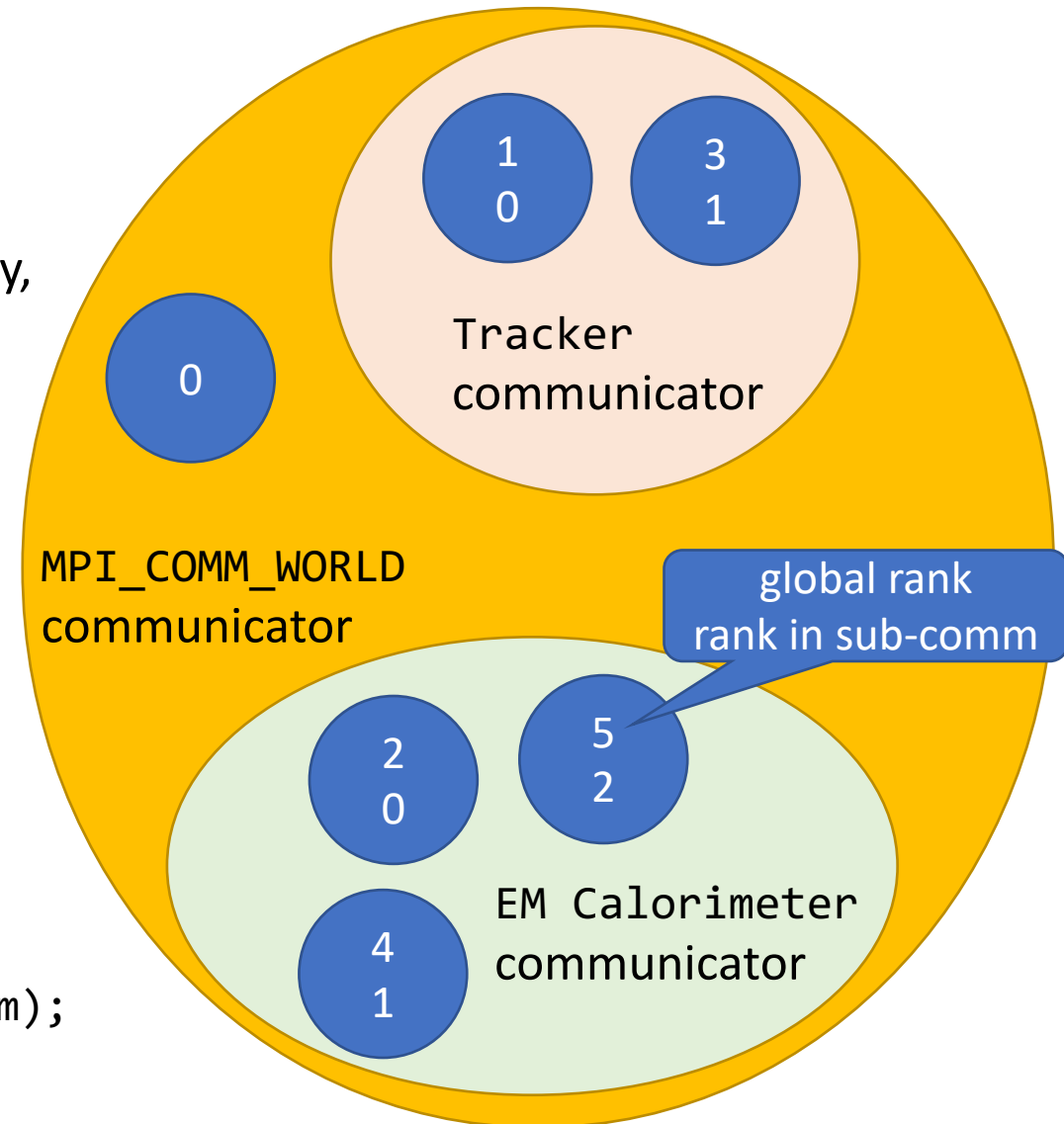
```
MPI_Comm newcomm;
```

```
MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);
```

```
mycolor = my_rank/3;
```

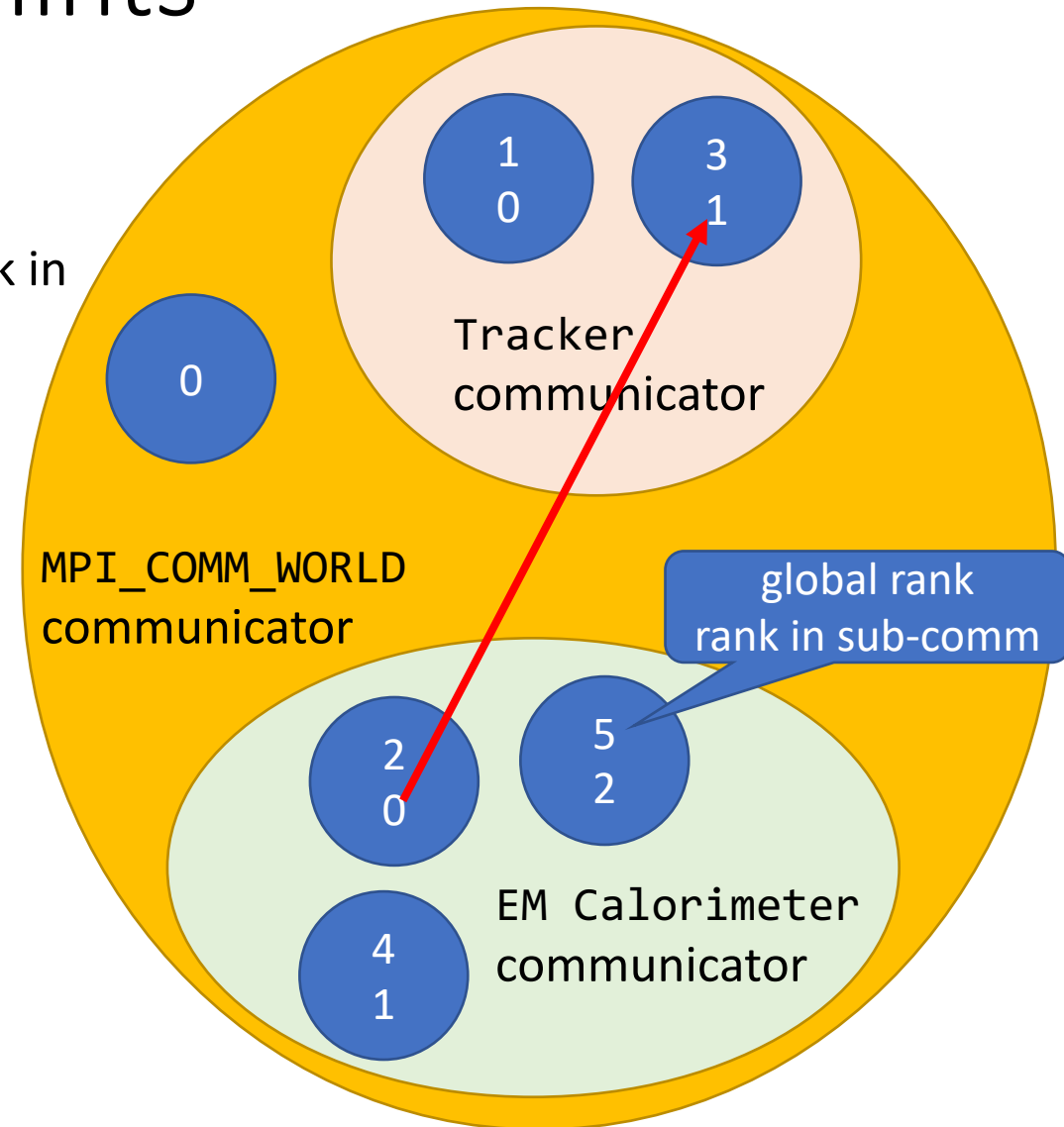
```
MPI_Comm_split(MPI_COMM_WORLD, mycolor, 0, &newcomm);
```

```
MPI_Comm_rank(newcomm, &my_new_rank);
```



MPI: Sub-Communicators hints

- no difference in speed: same hardware
- Use intra-communicators to communicate between rank in different “worlds” (without MPI_COMM_WORLD ranks)
 - eg MPI_Intercomm_create()





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Motivation: Collective Communication

- eg matrix multiplication, helpful:
 - reading and spreading of data,
 - gather final results

a_{00}	a_{01}	\cdots	$a_{0,n-1}$
a_{10}	a_{11}	\cdots	$a_{1,n-1}$
\vdots	\vdots		\vdots
a_{i0}	a_{i1}	\cdots	$a_{i,n-1}$
\vdots	\vdots		\vdots
$a_{m-1,0}$	$a_{m-1,1}$	\cdots	$a_{m-1,n-1}$

x_0
x_1
\vdots
x_{n-1}

 $=$

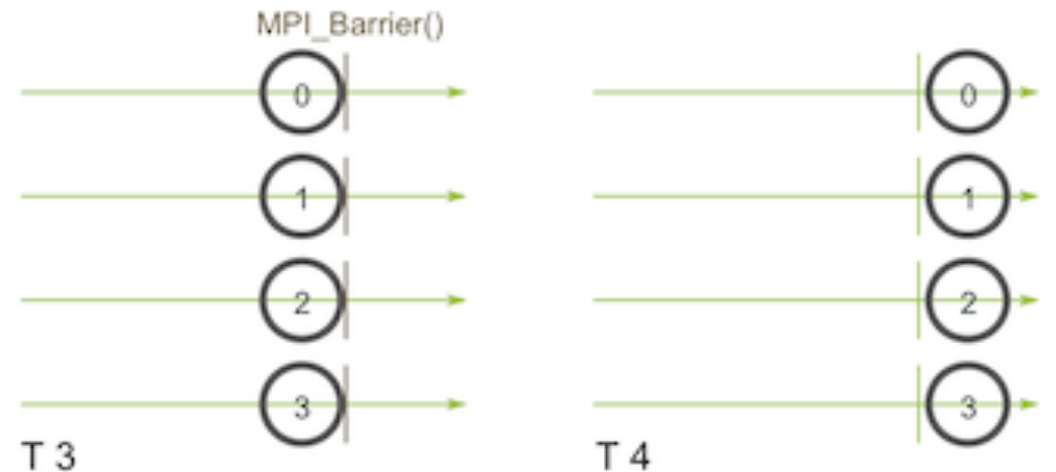
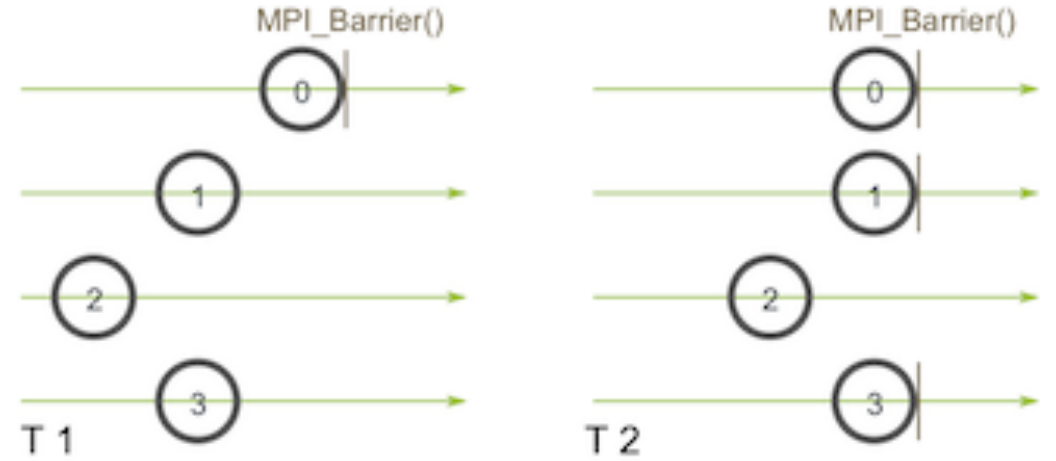
y_0
y_1
\vdots
$y_i = a_{i0}x_0 + a_{i1}x_1 + \cdots a_{i,n-1}x_{n-1}$
\vdots
y_{m-1}

MPI: MPI_Barrier

- collective communication: always include a *synchronization point* among processes.
 - all processes must reach a point in their code before they can all begin executing again.

syntax:

```
MPI_Barrier(MPI_Comm comm);
```

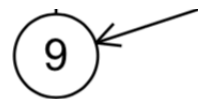
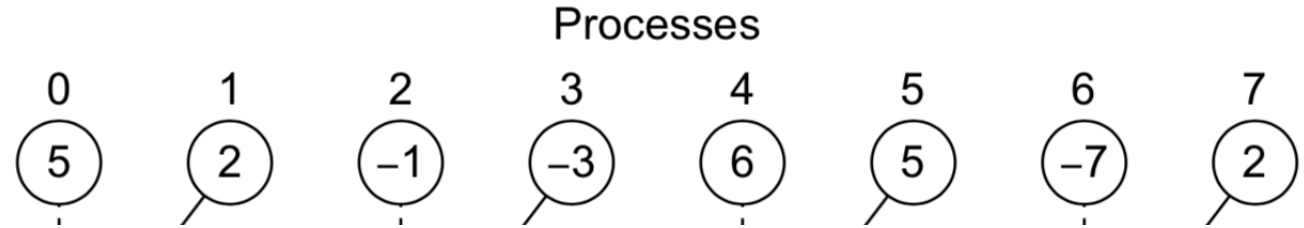


MPI: MPI_Reduce

- Reduces values on all processes to a single value (eg global sum)

```
int MPI_Reduce(  
void *sendbuf /*in*/,  
void *recvbuf /*out*/,  
int count /*in*/,  
MPI_Datatype datatype /*in*/,  
MPI_Op operator /*in*/,  
int dest_process /*in*/,  
MPI_Comm comm /*in*/)
```

- hints:
 - with count>1, MPI can operate on arrays
 - sendbuf and recvbuf need to be different (no aliasing!)

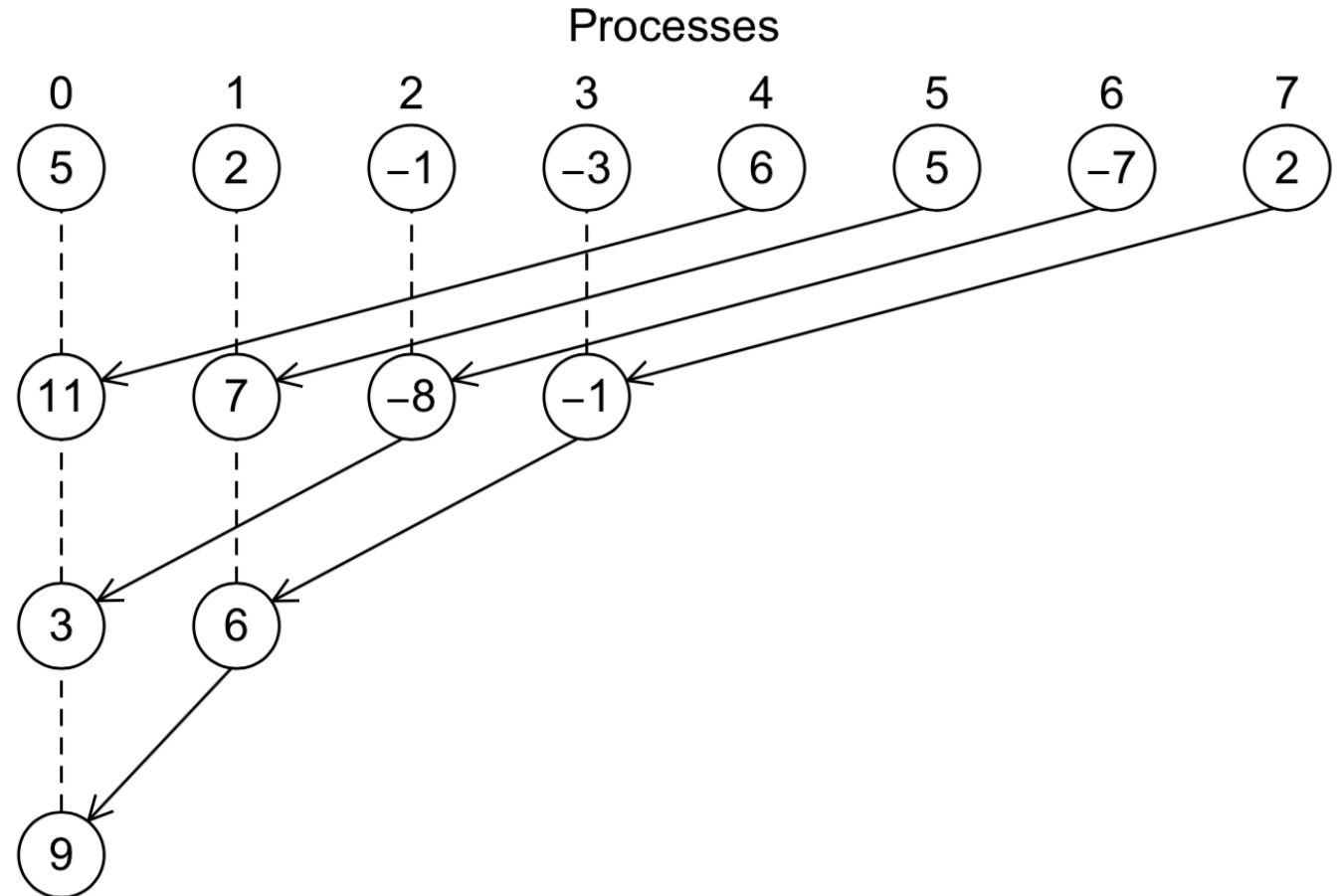


MPI: MPI_Reduce

- $\text{steps} = \lceil \log_2(N) \rceil$

```
int MPI_Reduce(  
void *sendbuf /*in*/,  
void *recvbuf /*out*/,  
int count /*in*/,  
MPI_Datatype datatype /*in*/,  
MPI_Op operator /*in*/,  
int dest_process /*in*/,  
MPI_Comm comm /*in*/)
```

- hint:
 - with $\text{count} > 1$, MPI can operate on arrays
 - sendbuf and recvbuf need to be different (no aliasing!)



MPI: MPI_Reduce

Worked out example:

```
int local_n, n;
```

```
local_n = my_rank;
```

```
MPI_Reduce(&local_n /*send_buf*/, &n /*recv_buf*/, 1 /*count*/, MPI_INT,  
          MPI_SUM, 0 /*dest_process*/, MPI_COMM_WORLD);
```

```
printf("sum of all local_n: %f", n);
```

MPI: Reduction Operators

Operation	Meaning
MPI_MAX	Returns the maximum element.
MPI_MIN	Returns the minimum element.
MPI_SUM	Sums the elements.
MPI_PROD	Multiplies all elements.
MPI_LAND	Performs a logical and across the elements.
MPI_LOR	Performs a logical or across the elements.
MPI_BAND	Performs a bitwise and across the bits of the elements.
MPI_BOR	Performs a bitwise or across the bits of the elements.
MPI_MAXLOC	Returns the maximum value and the rank of the process that owns it.
MPI_MINLOC	Returns the minimum value and the rank of the process that owns it.

MPI: P2P ↔ Collective Communication

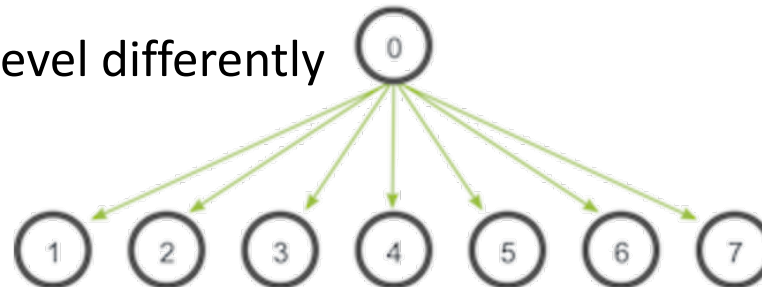
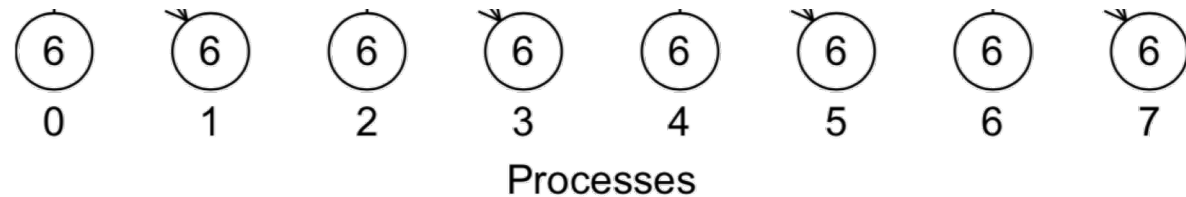
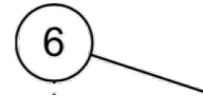
- ALL processes in communicator must call SAME collective function at the same time.
- Arguments in all ranks must fit:
 - eg. same dest_process, datatype, operator, comm
 - depending on function
- Only rank dest_process may use recvbuf (but all ranks have to provide such argument)
- MPI_Reduce calls matched solely on:
 - the communicator and
 - the order on which they are called.
 - No helping tags or sender id available.

MPI: Broadcast

Broadcasts the same message from the process "sending_rank" to all other processes of the communicator

- `MPI_Bcast`(
void *data,
int count,
MPI_Datatype datatype,
int sending_rank,
MPI_Comm comm)

- Hint: All ranks have to call this function
- Might be implemented on hardware level differently (MPI implementation should know)



MPI: Broadcast

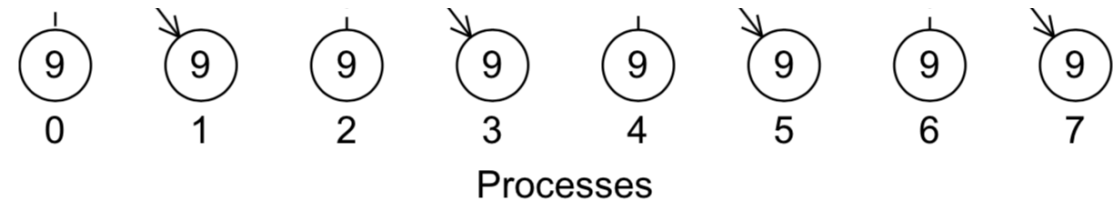
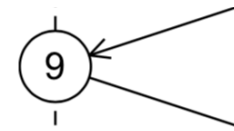
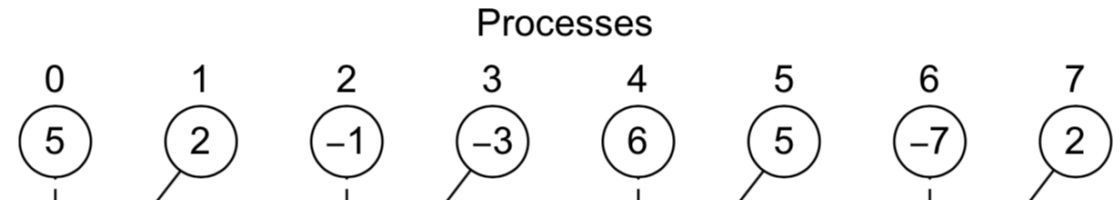
Worked out example, sending 2 variables:

```
int my_rank, my_size, n;
double a;

if (!my_rank) {
    printf("Enter a and n:\n");
    scanf("%lf %d", &a, &n);
}
MPI_Bcast(&a, 1, MPI_DOUBLE, 0, MPI_COMM_WORLD);
MPI_Bcast(&n, 1, MPI_DOUBLE, 0, MPI_COMM_WORLD);
```


MPI: MPI_Allreduce

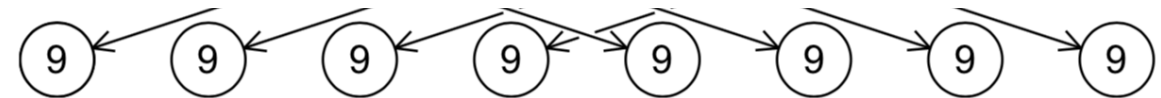
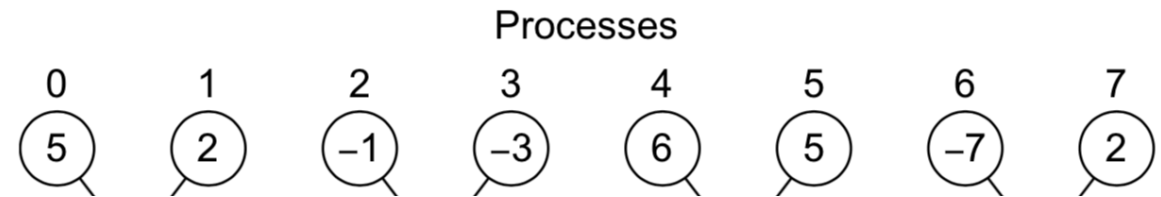
- Combines values from all processes and distributes the result back to all processes
eg: all processes need global sum
 1. compute global sum (MPI_Reduce) + Broadcast



MPI: MPI_Allreduce

- Combines values from all processes and distributes the result back to all processes
eg: all processes need global sum
 1. compute global sum (MPI_Reduce) + Broadcast
 2. exchange partial sums (better!, “butterfly”)
- MPI has optimal performance with

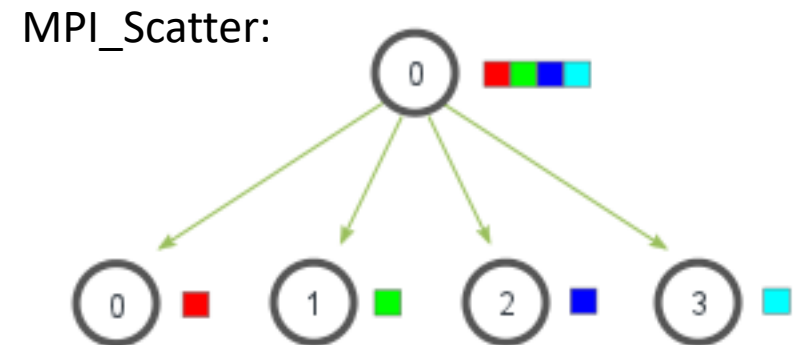
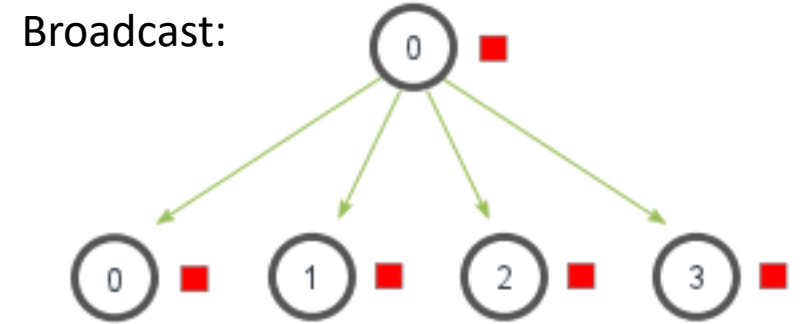
```
int MPI_Allreduce(  
void *sendbuf /*in*/,  
void *recvbuf /*out*/,  
int count /*in*/,  
MPI_Datatype datatype /*in*/,  
MPI_Op operator /*in*/,  
MPI_Comm comm /*in*/)
```



MPI: MPI_Scatter

Sends data from one process to all other processes in a communicator

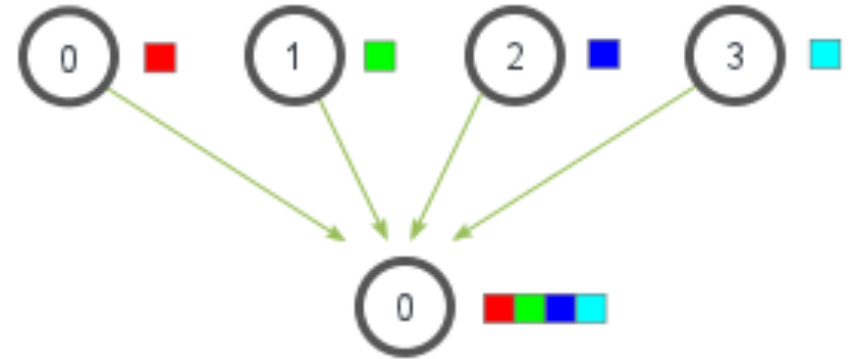
- Selective distribution of data to processes
- MPI_Scatter(
void* send_data /*in*/,
int send_count /*in*/,
MPI_Datatype send_datatype /*in*/,
void* recv_data /*out*/,
int recv_count /*in*/,
MPI_Datatype recv_datatype /*in*/,
int src_proc /*in*/,
MPI_Comm comm /*in*/)
- Special cases: MPI_Scatterv



MPI: MPI_Gather

Gathers together values from a group of processes

- MPI_Gather(
void* send_data /*in*/,
int send_count /*in*/,
MPI_Datatype send_datatype /*in*/,
void* recv_data /*out*/,
int recv_count /*in*/,
MPI_Datatype recv_datatype /*in*/,
int dest_proc /*in*/,
MPI_Comm comm /*in*/)
- Special cases: MPI_Gatherv



Exercises 5 and 6

- See online:
 - <https://gitlab.rlp.net/pbotte/learnhpc/tree/master/mpi>